# Mathematical Optimization for Social Distancing

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Picture from https://www.ted.com/playlists/735/the\_fight\_against\_viruses

Based on joint work with Martina Fischetti and Jakob Soustroup



# Facility location under social distancing constraints

The spread of viruses such as SARS-CoV-2 brought new challenges to our society, including a stronger focus on safety across all businesses





In particular, many countries have imposed a **minimum social distance** between people (or, more generally, **facilities**) to ensure their safety

## Can **Operations Research** help finding **more efficient** yet **safer** location patterns for facilities?

#### **Example: outdoor table allocation**



A brewpub in Copenhagen



Optimized layout fitting 6 more tables





Still 30 tables, but in globally safer configuration

### The safest distance problem

#### Given:

- An available (possibly disconnected/irregular) area
- A minimum and maximum n. of facilities to locate in that area
- A minimum distance between facilities
- A virus-spread model to measure infection probability among facilities
  - Find:

A facility location pattern that minimizes the overall infection probability (sum of all pairwise probabilities)







### **Applications**

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#### 2 variants of the problem:

 Fit as many customers as possible while respecting social distancing → a familiar packing problem



 Fit a fixed number of customers maximizing their safety → a new (?) quadratic optimization problem



### A mathematical optimization model

The available area is sampled to get a set
 V = {1,...,n} of possible allocation points



- Let  $N_{MIN}$  and  $N_{MAX}$  be min and max n. of facilities to allocate, resp.
- Let  $P_i$  denote the **profit** of allocating a facility at point *i* (e.g.,  $P_i \equiv 1$ )
- Define an incompatibity graph G<sub>I</sub>= (V, E<sub>I</sub>) whose edges [i,j] correspond to infeasible pairs with distance(i,j) < minimum\_distance</li>
- Let *I<sub>ij</sub>* denote the infection probability that *j* is infected by *i* (assuming *i* is positive)



#### The safest distance problem

$$\max\sum_{i\in V} P_i x_i - \sum_{i\in V} \sum_{j\in V} I_{ij} x_i x_j \tag{1}$$

s.t. 
$$N_{MIN} \le \sum_{i \in V} x_i \le N_{MAX}$$
 (2)

$$x_i + x_j \le 1 \qquad \qquad \forall [i, j] \in E_I \tag{3}$$

$$x_i \in \{0, 1\} \qquad \qquad \forall i \in V \tag{4}$$

#### An analogy: wind turbine allocation









#### An analogy: wind turbine allocation



## Wind farm layout optimization

#### Given

- a site (offshore or onshore)
- characteristics of the turbines to build
- measurements of the wind in the site



Determine a turbine allocation that maximizes power production

Taking into account:

- proximity constraints (no collisions)
- minimum/maximum number of turbines
- paiwise interference due to wake effects



## The wind farm layout problem

- Define a grid of candidate **positions** for turbine allocation
- For each pair of positions *i,j*, let *I<sub>ij</sub>* denote the average **interference** (power loss) experienced at position *j* if a turbine is built in position *i* →



it depends on average wind speed/direction, turbine power curve, etc.







only one wind scenario

three wind scenarios

several wind scenarios

• Assume overall interference is cumulative (sum of pairwise interf.s)

### **Basic (noncovex) quadratic MIP**

$$\max \sum_{i \in V} P_i x_i - \sum_{i \in V} \sum_{j \in V} I_{ij} x_i x_j$$
(1)  
s.t. 
$$N_{MIN} \leq \sum_{i \in V} x_i \leq N_{MAX}$$
(2)  
$$x_i + x_j \leq 1 \qquad \forall [i, j] \in E_I$$
(3)  
$$x_i \in \{0, 1\} \qquad \forall i \in V$$
(4)

#### **A standard MILP reformulation**

• Introduce a quadratic n. of var.s  $z_{ij} = x_i x_j$ 

$$\max \sum_{i \in V} P_i x_i - \sum_{i \in V} \sum_{j \in V, i < j} (I_{ij} + I_{ji}) z_{ij}$$
(1)

s.t. 
$$N_{MIN} \le \sum_{i \in V} x_i \le N_{MAX}$$
 (2)

$$x_i + x_j \le 1 \qquad \qquad \forall [i, j] \in E_I \tag{3}$$

$$x_i + x_j - 1 \le z_{ij} \qquad \forall i, j \in V, i < j \tag{4}$$

$$x_i \in \{0, 1\} \qquad \qquad \forall i \in V \tag{5}$$

$$z_{ij} \in \{0, 1\} \qquad \forall i, j \in V, i < j \tag{6}$$

#### **An alternative MILP reformulation**

Glover's trick: the objective function

$$\sum_{i \in V} P_i x_i - \sum_{i \in V} (\sum_{j \in V} I_{ij} x_j) x_i$$
(11)

is restated as

$$\sum_{i \in V} (P_i x_i - w_i) \tag{12}$$

where

$$w_i := \left(\sum_{j \in V} I_{ij} x_j\right) x_i = \begin{cases} \sum_{j \in V} I_{ij} x_j & \text{if } x_i = 1; \\ 0 & \text{if } x_i = 0. \end{cases}$$

→ the new continuous variable  $w_i$  is the product between a continuous term ( $\sum ...$ ) and a binary variable ( $x_i$ )

#### **An alternative MILP reformuation**

A linearized model with linear n. of additional var.s  $w_i$  and BIGM constr.s

$$\max z = \sum_{i \in V} (P_i x_i - w_i) \tag{13}$$

s.t. 
$$N_{MIN} \le \sum_{i \in V} x_i \le N_{MAX}$$
 (14)

$$x_i + x_j \le 1 \qquad \qquad \forall [i, j] \in E_I \qquad (15)$$

$$\sum_{j \in V} I_{ij} x_j \le w_i + M_i (1 - x_i) \quad \forall i \in V$$

$$x_i \in \{0, 1\} \quad \forall i \in V \quad (17)$$

$$w_i \ge 0 \quad \forall i \in V \quad (18)$$

where  $M_i >> 0$  (BIGM)

#### **Solution methods**

- No interference → packing problem → exactly solvable by modern MILP solvers up to hundreds of possible positions
- With interference  $\rightarrow$  very hard quadratic problem
- Proximity search matheuristics from the wind-farm opt. literature



Martina Fischetti and Michele Monaci (2016), Proximity search heuristics for wind farm optimal layout, *Journal of Heuristics* 22 (4), pp. 459-474.

 Results in the present talk have been obtained by an ad-hoc wind-farm layout optimization software developed by *Double-Click SRL*, Padua, Italy

## **Modelling virus spread**





Reliable wind turbine interference models in the literature (e.g., Jensen's model)

Virus spread models?

## **Modelling virus spread**

Let  $d_{ij}$  be the distance (in m) between points *i* and *j*, and let  $d_{max}$  be the max. distance in the given area

Alternative definitions for the virus-interference  $I_{ii}$ 

$$\begin{split} I_{ij} &= d_{max} - d_{ij} & \mathsf{P} \\ I_{ij} &= e^{-d_{ij}^2/2} & \mathsf{G} \\ I_{ij} &= 1/d_{ij} & \mathsf{U} \\ I_{ij} &= 1/d_{ij}^2 & \inf \\ I_{ij} &= 1/d_{ij}^3 & \mathsf{d} \end{split}$$

Proportional to distance

Gaussian ( $\sigma = 1m$ )

Uniform concentration inside a 1- or 2- or 3dimensional sphere



- Gaussian fits droplet (negligible after 3 σ)
- Uniform concentration fits aerosol spread



Example: optimal allocation of **10 facilities** on a line segment (no minimum distance imposed)

## **Modeling virus spread**



#### Optimal allocation of 20 facilities on a square

#### **Applications: a pub in Copenhagen**

![](_page_19_Figure_1.jpeg)

Minimum distance of 3m

![](_page_19_Figure_3.jpeg)

#### Regular solution

![](_page_19_Figure_5.jpeg)

#### **Applications: a pub in Copenhagen**

![](_page_20_Figure_1.jpeg)

Min dist. 3m but fixed n. of tables

![](_page_20_Figure_3.jpeg)

![](_page_20_Figure_4.jpeg)

#### **Applications: beach umbrellas in Venice**

![](_page_21_Picture_1.jpeg)

Min distance of 4m

8 12 16 20 24 28 32 36 40 44 48 52 56 60 64 68 72 76 80 84 88 

Optimized solution

8 more umbrellas

#### **Lessons learned**

- Facility location under social distance constraints and wind farm layout optimization are "similar" problems → we can borrow models and algorithms from wind farm literature
- Operations Research can make a big impact for businesses and customers, both in terms of profit and safety
- Fair layouts are often not regular and not easy to find "by hand"

![](_page_22_Picture_4.jpeg)

 Mathematical optimization is instrumental to finding more profitable yet safer configurations

![](_page_23_Picture_0.jpeg)

Full paper available at <a href="https://www.dei.unipd.it/~fisch/papers">www.dei.unipd.it/~fisch/papers</a>

Martina Fischetti, Matteo Fischetti, Jakob Stoustrup, Mathematical optimization for social distancing, Tech. Rep. UniPD, 2020.