# The consensus problem in the behavioral approach 

Mauro Bisiacco, Ingrid Blumthaler, and Maria Elena Valcher<br>Department of Information Engineering, University of Padova, Padova, Italy<br>\{bisiacco, meme\}@dei.unipd.it, ingrid.blumthaler@uibk.ac.at.

October 26, 2015

To Jan: a mentor and a friend. Amazingly inspiring and entertaining in both roles. $\mathscr{M} \mathscr{E} \mathscr{V}$

## 1 Introduction

The mathematical formulation of multi-agents systems and consensus problems was introduced several years ago in some pioneering papers such as [6, 17, 18]. But it was only a decade ago that a wide stream of research on these topics started, thanks to milestone contributions such as $[7,8,10,11,12]$. Aside from the theoretical challenges that these problems pose, strong motivations for such a widespread interest come from the numerous application problems that can be naturally stated as consensus problems. Indeed, when dealing with sensor networks, coordination of mobile robots or UAVs, flocking and swarming in animal groups, dynamics of opinion forming, etc., the main control target can be mathematically formalized as a consensus problem among agents, exchanging information and resorting to distributed algorithms that make use of the information collected from neighboring agents (see, e.g. [14, 16]).

While the first contributions on this subject focused on agents described as simple or double integrators, more recent works addressed the case of agents described by higher order models $[7,16,19,20,21]$. The vast majority of the literature on consensus, however, assumes that the homogeneous agents dynamics is described by a state-space model and that consensus is achieved through a static state- or output-feedback, that makes use of the weighted information collected from the neighboring agents, in a cooperative set-up (see [1] for consensus under antagonistic interactions).

The aim of this paper is to investigate the multi-agent consensus problem in a broader context, by assuming both for the agents and for the distributed controllers higher order input/output dynamic models. The behavioral approach developed by Jan Willems [13, 22, 23] seems to be a convenient set-up where to investigate this
general problem. Since, to the best of our knowledge, this set-up has never been used before in this context, we have tried to make the paper as self-contained as possible, by recalling the few fundamental definitions and results that are necessary to understand the technical details of the paper. A comprehensive treatment of the behavior theory can be found in any of the three aforementioned references.

By making use of the behavioral approach, we will show that the consensus problem can be naturally rephrased as a variant of the stabilization problem: the stabilization pertains only to a part of the system variables (the outputs) and it is achieved through regular full interconnection of the agents models and of the controllers. We will prove that if the communication among agents is described by a weighted, undirected and connected graph, then a necessary and sufficient condition for the consensus problem to be solvable is that the output is stabilizable from the input in the agents model. In this respect, the theory here developed for higher-order input/output models naturally extends the results about consensus derived in the state-space approach (see [20], for instance).

The paper is organized as follows. In section 2 preliminary definitions, notation and results are given. In section 3 the consensus problem is posed. Section 4 provides a characterization of the controllers that make it possible for the agents to achieve consensus and section 5 provides a similar characterization under the additional assumption that the consensus is achieved by means of a regular interconnection. Section 6 provides a complete solution to the consensus problem. Section 7 concludes the paper by showing how the most classical result on consensus for agents described by state-space models easily follows from the present analysis.

A preliminary version of the results appearing in sections 3,4 and 5 of this paper has appeablack, in a more general set-up, in [5]. However, no problem solution was provided: a characterization of the controllers that solve the problem was given, but no necessary and sufficient condition for the existence of such controllers, and hence for the consensus problem solvability, was provided.

## 2 Preliminaries

We introduce some notation and definitions that will be used in the following.
$I_{p}$ denotes the $p \times p$ identity matrix. The $p$-dimensional vector with all entries equal to 1 is denoted by $\mathbf{1}_{p}$, while the $i$ th standard basis vector in $\mathbb{R}^{p}$ (also known as the $i$ th canonical vector) is denoted by $e_{i}$. The spectrum of a square matrix $L$ is denoted by $\sigma(L)$. $\operatorname{diag}\left\{v_{1}, v_{2}, \ldots, v_{p}\right\}$ is the $p \times p$ diagonal matrix with diagonal entries $v_{1}, v_{2}, \ldots, v_{p}$.

We let $\mathbb{R}[s]$ denote the ring of polynomials in the indeterminate $s$ with real coefficients. A polynomial $p \in \mathbb{R}[s]$ is Hurwitz if all its zeros belong to $\{s \in \mathbb{C}$ : $\operatorname{Re}(s)<0\}$. A polynomial matrix $P=P(s) \in \mathbb{R}[s]^{p \times q}$ is right prime if it is of full column rank $q$ and the greatest common divisor of its maximal order minors is a unit, equivalently if $\operatorname{rank} P(\lambda)=q$ for every $\lambda \in \mathbb{C}$. It is well-known [9] that $P(s)$
is right prime if and only if it admits a polynomial left inverse or, equivalently, the Bézout equation

$$
X P=I_{q}
$$

in the unknown polynomial matrix $X(s) \in \mathbb{R}[s]^{q \times p}$ is solvable. Left prime matrices are similarly defined and characterized. A square and nonsingular polynomial matrix $P=P(s) \in \mathbb{R}[s]^{q \times q}$ whose inverse $P^{-1}$ is polynomial is called unimodular. Clearly a unimodular matrix is both right prime and left prime.

Every polynomial matrix $P \in \mathbb{R}[s]^{p \times q}$ of rank $r$ factorizes over $\mathbb{R}[s]$ as $P=L \Delta R$, where $L$ is $p \times r$ and right prime, $\Delta$ is $r \times r$ and nonsingular, and $R$ is $r \times q$ and left prime.

The concepts of left annihilator and, in particular, of minimal left annihilator (MLA, for short) of a given polynomial matrix $P$ have been originally introduced in [15] and can be summarized as follows: if $P$ is a $p \times q$ polynomial matrix of rank $r$, a polynomial matrix $H$ is a left annihilator of $P$ if $H P=0$. A left annihilator $H_{m}$ of $P$ is an MLA if it is of full row rank and for any other left annihilator $H$ of $P$ we have $H=Q H_{m}$ for some polynomial matrix $Q$. It can be easily proved that, unless $P$ is of full row rank, an MLA always exists (if $P$ is of full row rank, its left annihilators are zero matrices with an arbitrary number of rows), it is a $(p-r) \times p$ left prime matrix and is uniquely determined modulo a unimodular left factor.

In the paper we consider (continuous-time) signals defined on the time set $\mathbb{R}$. Signals will be real valued and hence they will be, in general, elements of $\left(\mathbb{R}^{q}\right)^{\mathbb{R}}$, for some $q \in \mathbb{N}$. By $\mathscr{F}^{q}$ we will denote the set of arbitrarily often differentiable functions, i.e., $\mathscr{C}^{\infty}\left(\mathbb{R}, \mathbb{R}^{q}\right) \subseteq\left(\mathbb{R}^{q}\right)^{\mathbb{R}}$.

For every $P=\sum_{i=0}^{n} P_{i} s^{i} \in \mathbb{R}[s]^{p \times q}$, we associate with $P$ the polynomial matrix differential operator $P \circ=\sum_{i=0}^{n} P_{i} \frac{d^{i}}{d t}$. The action of such a polynomial matrix differential operator $P$ on any signal $w \in \mathscr{F}^{q}$ is denoted by $P \circ w$.

In this paper by a system we mean a triple $\Sigma=\left(\mathbb{R}, \mathbb{R}^{q}, \mathscr{B}\right)$, where $\mathbb{R}$ is the time set, $\mathbb{R}^{q}$ is the set where the system trajectories take values, and $\mathscr{B}$ is the behavior, namely the set of admissible trajectories of the system variable $w$. We will consider linear, time-invariant behaviors described as the kernels of polynomial matrix operators. This means that there exists a polynomial matrix $P \in \mathbb{R}[s]^{k \times q}$ such that

$$
\begin{equation*}
\mathscr{B}=\left\{w \in \mathscr{F}^{q}: P \circ w=0\right\} . \tag{1}
\end{equation*}
$$

It is always possible to find a matrix $\bar{P} \in \mathbb{R}[s]^{r \times q}$ of full row rank $r$ such that $\mathscr{B}=$ $\left\{w \in \mathscr{F}^{q}: \bar{P} \circ w=0\right\}$.

A behavior $\mathscr{B} \subseteq \mathscr{F}^{q}$ is autonomous if it is a finite dimensional vector subspace of $\mathscr{F}^{q}$ as a vector space on $\mathbb{R} . \mathscr{B}$ described as in (1) is autonomous if and only if $P \in \mathbb{R}[s]^{k \times q}$ is of full column rank $q$.

An autonomous behavior (1) is stable if the greatest common divisor of the maximal (i.e., $q$ th) order minors of $P$ is a Hurwitz polynomial. If $P$ is of full row rank and hence, under the autonomy assumption, square and nonsingular, this amounts to requiring that $\operatorname{det} P$ is Hurwitz. A trajectory $w \in \mathscr{F}^{q}$ is called small
if it belongs to some stable autonomous behavior or, equivalently, if it satisfies the equation $p \circ w=0$ for some Hurwitz polynomial $p$. Clearly, small signals are the polynomial-exponential functions that converge to zero as the time approaches $+\infty$.

If we partition the system variables as $w=\binom{y}{u} \in \mathscr{F}^{p+m}$, and accordingly describe the behavior $\mathscr{B}$ as

$$
\mathscr{B}=\left\{\binom{y}{u} \in \mathscr{F}^{p+m}: P_{y} \circ y=P_{u} \circ u\right\}, \quad\left(P_{y}-P_{u}\right) \in \mathbb{R}[s]^{k \times(p+m)},
$$

we say that $u$ is free in $\mathscr{B}$ if for any $u \in \mathscr{F}^{m}$ there exists $y \in \mathscr{F}^{p}$ such that $\binom{y}{u} \in \mathscr{B}$. This is the case if and only if $\operatorname{rank}\left(P_{y}-P_{u}\right)=\operatorname{rank}\left(P_{y}\right)$. If additionally the behavior

$$
\mathscr{B}^{0}=\left\{y \in \mathscr{F}^{p}: P_{y} \circ y=0\right\}
$$

is autonomous, we say that $\mathscr{B}$ is an input/output behavior with input $u$ and output $y$. Clearly, this is the case if and only if $\operatorname{rank}\left(P_{y}-P_{u}\right)=\operatorname{rank}\left(P_{y}\right)=p$. If the matrix $\left(P_{y}-P_{u}\right)$ is of full row rank $k$, it follows that $\mathscr{B}$ is an input/output behavior with input $u$ and output $y$ if and only if $k=p$ and $P_{y}$ is nonsingular.

If $\mathscr{B}$ and $\mathscr{C}$ are behaviors in $\mathscr{F}^{q}$, described as kernels of the polynomial matrix operators $P \circ$ and $C \circ$, respectively, we denote the interconnection of $\mathscr{B}$ and $\mathscr{C}$ as follows:

$$
\mathscr{B} \wedge \mathscr{C}:=\left\{w \in \mathscr{F}^{q}: w \in \mathscr{B}, w \in \mathscr{C}\right\}=\left\{w \in \mathscr{F}^{q}:\binom{P}{C} \circ w\right\} .
$$

The interconnection of $\mathscr{B}$ and $\mathscr{C}$ is said to be regular if

$$
\operatorname{rank}\binom{P}{C}=\operatorname{rank}(P)+\operatorname{rank}(C)
$$

If $\mathscr{B}$ is an input/output behavior, with input $u$ and output $y$, and the interconnection of $\mathscr{B}$ and $\mathscr{C}$ is regular, the input/output structure of $\mathscr{B}$ is preserved even after interconnection: this means that it is still possible to add (free) signals $u^{\prime}$ to the components of $u$ after interconnection [2]. More precisely, the components of $u^{\prime}$ are free in the behavior $\mathscr{B} \wedge \mathscr{C}$, illustrated in Figure 1.

The concept of regular interconnection is fundamental in the behavioral approach and will be used in the sequel. For an in-depth discussion of its meaning and relevance see $[2,24]$.

Given a (not necessarily an input/output) behavior $\mathscr{B}$, with signals $\binom{y}{u} \in \mathscr{F}^{p+m}$, we say that $y$ is stabilizable from $u$ in $\mathscr{B}$ [2] if for every pair of trajectories $(y, u) \in$ $\mathscr{B}$ there exists $(\bar{y}, \bar{u}) \in \mathscr{B}$ such that $\bar{y}(t)=y(t), \forall t<0$, and $\bar{y}$ converges to zero. By rephrasing the result obtained in [3] for zero-controllability of discrete-time systems, we obtain the following characterization.


Figure 1: Feedback interconnection of the input/output behavior $\mathscr{B}$ and $\mathscr{C}$, with additional additive inputs $u^{\prime}$.

Proposition 1. Given a behavior

$$
\mathscr{B}=\left\{\binom{y}{u} \in \mathscr{F}^{p+m}: P_{y} \circ y=P_{u} \circ u\right\}
$$

the following facts are equivalent:
i) $y$ is stabilizable from $u$ in $\mathscr{B}$;
ii) there exists a behavior $\mathscr{C}$ in $\mathscr{F}^{p+m}$ such that the interconnection $\mathscr{B} \wedge \mathscr{C}$ is regular, and the projection of $\mathscr{B} \wedge \mathscr{C}$ on the variable y, i.e., $\left\{y \in \mathscr{F}^{p}:(y, u) \in\right.$ $\left.\mathscr{B} \wedge \mathscr{C}, \exists u \in \mathscr{F}^{m}\right\}$, is an autonomous and stable behavior;
iii) either one of the following two cases applies:

1. $P_{u}$ is of full row rank, or
2. $P_{u}$ is not of full row rank, and if we let $M_{u}$ denote an MLA for $P_{u}$, then $M_{u} P_{y}$ factorizes over $\mathbb{R}[s]$ as

$$
M_{u} P_{y}=L \Delta R
$$

where $L$ is right prime, $\Delta$ is square and Hurwitz, and $R$ is left prime.
A behavior $\mathscr{C}$ such that the interconnected system $\mathscr{B} \wedge \mathscr{C}$ is an autonomous stable behavior is called a stabilizing controller for $\mathscr{B}$ [2, 13]. Condition ii) in Proposition 1 means something weaker, namely that the controller $\mathscr{C}$ is stabilizing only for the variable $y$ of the behavior $\mathscr{B}$.

## 3 The consensus problem

We assume that there are $N \geq 2$ agents whose dynamics is described by the same input/output behavior

$$
\begin{align*}
\mathscr{P}_{i}:=\mathscr{P}:= & \left\{\binom{y_{i}}{u_{i}} \in \mathscr{F}^{p+m}: P_{y} \circ y_{i}=P_{u} \circ u_{i}\right\},  \tag{2}\\
& \text { with }\left(P_{y}-P_{u}\right) \in \mathbb{R}[s]^{p \times(p+m)}, \quad \operatorname{det} P_{y} \neq 0,
\end{align*}
$$

for $i=1, \ldots, N$. The variables $y_{i}$ represent the target variables, i.e., those on which the agents should reach consensus. We assume that they are measurable (and hence available for feedback interconnection). Note that the assumption $\operatorname{det}\left(P_{y}\right) \neq 0$ ensures that the matrix $\left(P_{y}-P_{u}\right)$ is of full row rank $p$.

Upon setting

$$
y:=\left(\begin{array}{c}
y_{1} \\
\vdots \\
y_{N}
\end{array}\right) \in \mathscr{F}^{N p}, \quad \text { and } \quad u:=\left(\begin{array}{c}
u_{1} \\
\vdots \\
u_{N}
\end{array}\right) \in \mathscr{F}^{N m},
$$

the overall behavior of the $N$ agents can be described as follows:

$$
\begin{aligned}
\mathscr{P} & :=\left\{\binom{y}{u} \in \mathscr{F}^{N p+N m}:\binom{y_{i}}{u_{i}} \in \mathscr{P}_{i}, \forall i \in\{1, \ldots, N\}\right\} \\
& =\left\{\binom{y}{u} \in \mathscr{F}^{N p+N m}:\left(I_{N} \otimes P_{y}\right) \circ y=\left(I_{N} \otimes P_{u}\right) \circ u\right\} \\
& \cong \mathscr{P}_{1} \times \ldots \times \mathscr{P}_{N},
\end{aligned}
$$

where $\otimes$ denotes the Kronecker product of matrices, and $\cong$ means that $\mathscr{P}$ and $\mathscr{P}_{1} \times \ldots \times \mathscr{P}_{N}$ are isomorphic, since their trajectories are mutually related by a simple entry permutation.

The information flow between the agents is modeled by an undirected and connected weighted $\operatorname{graph} \mathscr{G}$, with vertex set $\mathscr{V}=\{1, \ldots, N\}$ (the $i$ th vertex representing the $i$ th agent) and edge set $\mathscr{E} \subseteq \mathscr{V} \times \mathscr{V}$. The adjacency matrix of $\mathscr{G}$ is nonnegative, symmetric and denoted by $A=\left(a_{i j}\right)_{1 \leq i, j \leq N} \in \mathbb{R}_{+}^{N \times N}: a_{i j}$ is the (positive) weight of the edge from $j$ to $i$ if such an edge exists, and zero otherwise. We assume that $a_{i i}=0$ for every $i \in\{1,2, \ldots, N\}$. Note that the assumption that the graph is undirected corresponds to the rather common case when the $i$ th agent gives to the information received from the $j$ th agent the same weight that the $j$ th agent gives to the information received from the $i$ th agent. The adjacency matrix gives rise to the Laplacian matrix $L:=\Delta-A$ with $\Delta:=\operatorname{diag}\left(A \mathbf{1}_{N}\right)$. The Laplacian matrix $L$ is a symmetric and positive semidefinite matrix, whose nonnegative real eigenvalues can be sorted in such a way that

$$
0 \leq \lambda_{1} \leq \lambda_{2} \leq \cdots \leq \lambda_{N}
$$

The connectedness assumption on $\mathscr{G}$ ensures that 0 is a simple eigenvalue of $L$, and hence all the eigenvalues $\lambda_{i}, i=2,3, \ldots, N$ are positive. Also, by the way $L$
has been defined, $\mathbf{1}_{N}$ is an eigenvector of $L$ corresponding to the zero eigenvalue. As $L$ is symmetric and hence diagonalizable, there exist $N$ linearly independent eigenvectors $v_{i}, i=1,2, \ldots, N$, with $v_{i}$ eigenvector corresponding to $\lambda_{i}$, such that the matrix

$$
V=\left(\begin{array}{llll}
v_{1} & v_{2} & \ldots & v_{N} \tag{3}
\end{array}\right)
$$

satisfies $L V=V D$, where $D$ is the diagonal matrix $\operatorname{diag}\left\{0, \lambda_{2}, \ldots, \lambda_{N}\right\}$. As $v_{1}$ is an eigenvector of $L$ corresponding to the zero eigenvalue, it entails no loss of generality assuming that it coincides with $\mathbf{1}_{N}$.

Each $i$ th agent receives the weighted information

$$
\widetilde{y}_{i}:=\sum_{j=1}^{N} a_{i j}\left(y_{i}-y_{j}\right)
$$

from the other agents of the network, and designs a control strategy based on the signal $\widetilde{y}_{i}$. Specifically, we assume that the agents adopt (identical) input/output controllers (with input $\widetilde{y}_{i}$ and output $u_{i}$ ) described by

$$
\begin{gather*}
\mathscr{C}_{i}:=\mathscr{C}:=\quad\left\{\binom{\widetilde{y}_{i}}{u_{i}} \in \mathscr{F} \mathscr{F}^{+m}: C_{y} \circ \widetilde{y}_{i}=C_{u} \circ u_{i}\right\}  \tag{4}\\
\text { where } \quad\left(C_{y}-C_{u}\right) \in \mathbb{R}[s]^{m \times(p+m)}, \quad \operatorname{det}\left(C_{u}\right) \neq 0
\end{gather*}
$$

for $i=1, \ldots, N$, and they give rise to the overall controller

$$
\begin{aligned}
& \underline{C}: \\
&:\left\{\binom{\widetilde{y}}{u} \in \mathscr{F}^{N p+N m}:\binom{\tilde{y}_{i}}{u_{i}} \in \mathscr{C}_{i}, \forall i \in\{1, \ldots, N\}\right\} \\
&=\left\{\binom{\widetilde{y}}{u} \in \mathscr{F}^{N p+N m}:\left(I_{N} \otimes C_{y}\right) \circ \widetilde{y}=\left(I_{N} \otimes C_{u}\right) \circ u\right\} \\
& \cong \mathscr{C}_{1} \times \ldots \times \mathscr{C}_{N}
\end{aligned}
$$

where, as before, $\tilde{y}:=\left(\begin{array}{c}\widetilde{y}_{1} \\ \vdots \\ \tilde{y}_{N}\end{array}\right) \in \mathscr{F}^{N p}, u=\left(\begin{array}{c}u_{1} \\ \vdots \\ u_{N}\end{array}\right) \in \mathscr{F}^{N m}$.
The agents behaviors $\mathscr{P}_{1}, \ldots \mathscr{P}_{N}$, the graph $\mathscr{G}$, and the controllers $\mathscr{C}_{1}, \ldots, \mathscr{C}_{N}$ define the overall interconnected behavior

$$
\left.\left.\begin{array}{rl}
\mathscr{K}:= & \left\{\binom{y}{u} \in \mathscr{F}^{N p+N m}: \forall i \in\{1, \ldots, N\},\right. \\
& =\left\{\binom{y_{i}}{u_{i}} \in \mathscr{P}_{i},\binom{\widetilde{y}_{i}}{u_{i}} \in \mathscr{C}_{i}, \text { where } \widetilde{y}_{i}=\sum_{j=1}^{N} a_{i j}\left(y_{i}-y_{j}\right)\right\} \\
u
\end{array}\right) \in \mathscr{F}^{N p+N m}:\binom{y}{u} \in \underline{\mathscr{P}},\binom{\widetilde{y}}{u} \in \underline{\mathscr{C}}, \text { where } \widetilde{y}=\left(L \otimes I_{p}\right) y\right\},
$$

This overall interconnection can also be interpreted as the interconnection of the overall plant $\mathscr{P}$ with the compensator

$$
\begin{aligned}
\mathscr{C}_{L} & :=\left\{\binom{y}{u} \in \mathscr{F}^{N p+N m}:\binom{\left(L \otimes I_{p}\right) y}{u} \in \underline{\mathscr{C}}\right\} \\
& =\left\{\binom{y}{u} \in \mathscr{F}^{N p+N m}:\left(L \otimes C_{y}\right) \circ y=\left(I_{N} \otimes C_{u}\right) \circ u\right\},
\end{aligned}
$$

namely as

$$
\mathscr{K}=\underline{\mathscr{P}} \wedge \underline{\mathscr{C}}_{L} .
$$

Definition 2. We say that the controllers $\mathscr{C}_{i}, i=1, \ldots, N$, lead to consensus among the $N$ agents (described by $\mathscr{P}_{1}, \ldots, \mathscr{P}_{N}$ and the graph $\mathscr{G}$ ) if for every trajectory $\binom{y}{u} \in \mathscr{K}$ the deviations $y_{j}-y_{1}, j \in\{2, \ldots, N\}$, are small.

## 4 Characterization of controllers leading to consensus

In order to characterize the controllers that lead to consensus, it is convenient to introduce a variable transformation that transforms the overall connected system $\mathscr{K}$ into the following isomorphic behavior:

$$
\mathscr{K}_{D}:=\left\{\binom{\bar{y}}{\bar{u}} \in \mathscr{F}^{N p+N m}:\left(\begin{array}{ll}
I_{N} \otimes P_{y} & -I_{N} \otimes P_{u} \\
D \otimes C_{y} & -I_{N} \otimes C_{u}
\end{array}\right) \circ\binom{\bar{y}}{\bar{u}}=0\right\} .
$$

As before, we assume that $\bar{y}=:\left(\begin{array}{c}\bar{y}_{1} \\ \vdots \\ \overline{y_{N}}\end{array}\right) \in \mathscr{F}^{N p}$ and $\bar{u}=:\left(\begin{array}{c}\bar{u}_{1} \\ \vdots \\ \bar{u}_{N}\end{array}\right) \in \mathscr{F}^{N m}$.
The behavior $\mathscr{K}_{D}$ can be interpreted as the interconnection of the overall plant $\mathscr{P}$ with the compensator

$$
\underline{\mathscr{C}}_{D}:=\left\{\left.\binom{\bar{y}}{\bar{u}} \in \mathscr{F}^{N p+N m} \right\rvert\,\left(D \otimes C_{y}\right) \circ \bar{y}=\left(I_{N} \otimes C_{u}\right) \circ \bar{u}\right\},
$$

namely as

$$
\mathscr{K}_{D}=\underline{\mathscr{P}} \wedge \underline{\mathscr{C}}_{D} .
$$

Note that, as the adjacency matrix $A$ of $\mathscr{G}$ is symmetric, and hence the Laplacian matrix $L$ is diagonalizable, the behavior $\mathscr{K}_{D}$ is the direct product of $N$ independent behaviors (see Figure 2 for an illustration of the structures of $\mathscr{K}$ and $\mathscr{K}_{D}$ ).

As a matter of fact, we have

$$
\begin{aligned}
\mathscr{K}_{D} & =\left\{\binom{\bar{y}}{\bar{u}} \in \mathscr{F}^{N p+N m}:\binom{\bar{y}_{i}}{\bar{u}_{i}} \in \mathscr{P},\binom{\lambda_{i} \bar{y}_{i}}{\bar{u}_{i}} \in \mathscr{C}, \forall i \in\{1, \ldots, N\}\right\} \\
& =\left\{\binom{\bar{y}}{\bar{u}} \in \mathscr{F}^{N p+N m}:\binom{\bar{y}_{i}}{\bar{u}_{i}} \in \mathscr{P} \wedge \mathscr{C}\left(\lambda_{i}\right), \forall i \in\{1, \ldots, N\}\right\} \\
& \cong\left(\mathscr{P} \wedge \mathscr{C}\left(\lambda_{1}\right)\right) \times \ldots \times\left(\mathscr{P} \wedge \mathscr{C}\left(\lambda_{N}\right)\right),
\end{aligned}
$$

where

$$
\begin{equation*}
\mathscr{C}\left(\lambda_{i}\right):=\left\{\binom{\bar{y}_{i}}{\bar{u}_{i}} \in \mathscr{F}^{p+m}:\binom{\lambda_{i} \bar{y}_{i}}{\bar{u}_{i}} \in \mathscr{C}\right\} . \tag{5}
\end{equation*}
$$

Lemma 3. The behaviors $\mathscr{K}$ and $\mathscr{K}_{D}$ are related through the following isomorphism (variable transformation):

$$
\binom{\bar{y}}{\bar{u}}=\left(\begin{array}{cc}
V^{-1} \otimes I_{p} & 0  \tag{6}\\
0 & V^{-1} \otimes I_{m}
\end{array}\right)\binom{y}{u},
$$

where $V$ is the matrix of eigenvectors of $L$ defined in (3), with $v_{1}=\mathbf{1}_{N}$, and satsfying $L V=V D$.

Proof. As $V \in \mathbb{R}^{N \times N}$ is nonsingular, it follows that also

$$
\underline{V}:=\left(\begin{array}{cc}
V \otimes I_{p} & 0 \\
0 & V \otimes I_{m}
\end{array}\right)
$$

is invertible over $\mathbb{R}[s] . \underline{V}^{-1}$ obviously defines an $\mathbb{R}[s]$-isomorphism from $\mathscr{F}^{N p+N m}$ into $\mathscr{F}^{N p+N m}$. We show that $\binom{y}{u} \in \mathscr{K}$ if and only if $\left(\frac{\bar{y}}{u}\right) \in \mathscr{K}_{D}$. By definition, $\binom{y}{u} \in \mathscr{K}$ if and only if

$$
\underbrace{\left(\begin{array}{cc}
I_{N} \otimes P_{y} & -I_{N} \otimes P_{u} \\
L \otimes C_{y} & -I_{N} \otimes C_{u}
\end{array}\right)}_{=: K} \circ\binom{y}{u}=0
$$

or equivalently if and only if

$$
\left(\left(\begin{array}{cc}
V^{-1} \otimes I_{n+p} & 0 \\
0 & V^{-1} \otimes I_{m}
\end{array}\right) K \underline{V}\right) \circ \underline{V}^{-1}\binom{y}{u}=0
$$

namely

$$
\left(\begin{array}{ll}
I_{N} \otimes P_{y} & -I_{N} \otimes P_{u} \\
D \otimes C_{y} & -I_{N} \otimes C_{u}
\end{array}\right) \circ\binom{\bar{y}}{\bar{u}}=0
$$

i.e., if and only if $\left(\frac{\bar{y}}{u}\right) \in \mathscr{K}_{D}$.

In light of Lemma 3 and of the decoupled expression of $\mathscr{K}_{D}$, we can derive the following characterization of the controllers that allow to solve the consensus problem.

Theorem 4. The following statements are equivalent:

1. The $N$ controllers $\mathscr{C}_{i}=\mathscr{C}, i=1,2, \ldots, N$, lead to consensus, i.e.,

$$
\forall\binom{y}{u} \in \mathscr{K}=\underline{\mathscr{P}} \wedge \underline{\mathscr{C}}_{L}, \forall i \in\{2, \ldots, N\}, \quad y_{i}-y_{1} \quad \text { is small. }
$$



Figure 2: The interconnections leading to $\mathscr{K}$ (original interconnected system) and $\mathscr{K}_{D}$ (after variable transformation).
2. The behavior $\mathscr{K}_{D}$ satisfies

$$
\forall\binom{\bar{y}}{\bar{u}} \in \mathscr{K}_{D}=\underline{\mathscr{P}} \wedge \underline{\mathscr{C}}_{D}, \forall i \in\{2, \ldots, N\}, \quad \bar{y}_{i} \quad \text { is small. }
$$

3. $\forall i \in\{2, \ldots, N\}, \forall\binom{\bar{y}_{i}}{\bar{u}_{i}} \in \mathscr{P} \wedge \mathscr{C}\left(\lambda_{i}\right), \quad \bar{y}_{i} \quad$ is small.
4. $\forall i \in\{2, \ldots, N\}, \forall\binom{\bar{y}_{i}}{\hat{u}_{i}} \in \mathscr{P}\left(\lambda_{i}\right) \wedge \mathscr{C}, \quad \bar{y}_{i} \quad$ is small, i.e., $\mathscr{C}$ stabilizes the variable $\bar{y}_{i}$ in $\mathscr{P}\left(\lambda_{i}\right) \wedge \mathscr{C}$, where

$$
\begin{equation*}
\mathscr{P}\left(\lambda_{i}\right):=\left\{\binom{\bar{y}_{i}}{\hat{u}_{i}} \in \mathscr{F}^{p+m}:\binom{\bar{y}_{i}}{\lambda_{i} \hat{u}_{i}} \in \mathscr{P}\right\}, \tag{7}
\end{equation*}
$$

and $\lambda_{i}, i=2,3, \ldots, N$ are the positive eigenvalues of $L$.
Proof. We first prove the equivalence of 1. and 2. Upon introducing the $(N-1) \times$ $N$ matrix $S:=\left(\begin{array}{ll}-\mathbf{1}_{N-1} & I_{N-1}\end{array}\right)$, statement 1. can be rewritten as

$$
\forall\binom{y}{u} \in \mathscr{K}, \quad\left(S \otimes I_{n}\right) y \quad \text { is small. }
$$

On the other hand, by Lemma $3,\binom{y}{u} \in \mathscr{K}$ if and only if $\binom{\bar{y}}{\bar{u}} \in \mathscr{K}_{D}$, where $\bar{u}$ and $\bar{y}$ are related to $u$ and $y$ via equation (6). In particular, $y=\left(V \otimes I_{n}\right) \bar{y}$, where $V$ is the usual matrix of eigenvectors that diagonalizes the Laplacian $L$. So, condition 1. can be equivalently rewritten as

$$
\begin{equation*}
\forall\binom{\bar{y}}{\bar{u}} \in \mathscr{K}_{D}, \quad\left(S \otimes I_{n}\right)\left(V \otimes I_{n}\right) \bar{y}=\left(S V \otimes I_{n}\right) \bar{y} \quad \text { is small. } \tag{8}
\end{equation*}
$$

The fact that $V$ is invertible and the identity
$V=\left(\begin{array}{cccc}1 & V_{12} & \cdots & V_{1 N} \\ 1 & V_{22} & \cdots & V_{2 N} \\ \vdots & \vdots & & \vdots \\ 1 & V_{N 2} & \cdots & V_{N N}\end{array}\right)=\left(\begin{array}{cccc}1 & 0 & \cdots & 0 \\ 1 & 1 & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 1 & 0 & \cdots & 1\end{array}\right)\left(\begin{array}{cccc}1 & V_{12} & \cdots & V_{1 N} \\ 0 & V_{22}-V_{12} & \cdots & V_{2 N}-V_{1 N} \\ \vdots & \vdots & & \vdots \\ 0 & V_{N 2}-V_{12} & \cdots & V_{N N}-V_{1 N}\end{array}\right)$
imply that also

$$
\widetilde{V}:=\left(\begin{array}{ccc}
V_{22}-V_{12} & \cdots & V_{2 N}-V_{1 N} \\
\vdots & & \vdots \\
V_{N 2}-V_{12} & \cdots & V_{N N}-V_{1 N}
\end{array}\right) \in \mathbb{R}^{(N-1) \times(N-1)}
$$

is invertible over $\mathbb{R}$. Since $S V=(0 \widetilde{V}) \in \mathbb{R}^{(N-1) \times(1+(N-1))}$, condition (8) is equivalent to

$$
\forall\binom{\bar{y}}{\bar{u}} \in \mathscr{K}_{D}, \quad\left((0 \widetilde{V}) \otimes I_{n}\right) \bar{y} \quad \text { is small }
$$

and by the invertibility of $\widetilde{V}$ over $\mathbb{R}$, this is equivalent to

$$
\forall\binom{\bar{y}}{\bar{u}} \in \mathscr{K}_{D}, \quad\left(\left(0 I_{N-1}\right) \otimes I_{n}\right) \bar{y} \quad \text { is small, }
$$

i.e., statement 2. is satisfied.

The equivalence of 2 . and 3 . follows from the special form of $\mathscr{K}_{D}$ as $D$ is diagonal. So, we are left with proving that the statements 3 . and 4. are equivalent.

For every $i \in\{2, \ldots, N\}$, the following equivalences hold:

$$
\begin{aligned}
\binom{\bar{y}_{i}}{\hat{u}_{i}} \in \mathscr{P}\left(\lambda_{i}\right) \wedge \mathscr{C} & \Leftrightarrow\binom{\bar{y}_{i}}{\lambda_{i} \hat{u}_{i}} \in \mathscr{P},\binom{\bar{y}_{i}}{\hat{u}_{i}} \in \mathscr{C} \\
& \Leftrightarrow\binom{\bar{y}_{i}}{\lambda_{i} \hat{u}_{i}} \in \mathscr{P},\binom{\lambda_{i} \bar{y}_{i}}{\lambda_{i} \hat{u}_{i}} \in \mathscr{C} \\
& \Leftrightarrow\binom{\bar{y}_{i}}{\lambda_{i} \hat{u}_{i}} \in \mathscr{P},\binom{\bar{y}_{i}}{\lambda_{i} \hat{u}_{i}} \in \mathscr{C}\left(\lambda_{i}\right) \\
& \Leftrightarrow\binom{\bar{y}_{i}}{\lambda_{i} \hat{u}_{i}} \in \mathscr{P} \wedge \mathscr{C}\left(\lambda_{i}\right) .
\end{aligned}
$$

Consequently, $\forall\binom{\bar{y}_{i}}{\bar{u}_{i}} \in \mathscr{P} \wedge \mathscr{C}\left(\lambda_{i}\right), \bar{y}_{i}$ is a small trajectory if and only if, $\forall\binom{\bar{y}_{i}}{\hat{u}_{i}} \in$ $\mathscr{P}\left(\lambda_{i}\right) \wedge \mathscr{C}, \bar{y}_{i}$ is a small trajectory. This proves that 3 . is equivalent to 4.

The main contribution of the previous theorem lies in the equivalence of conditions 1. and 4. Indeed, it shows that the $N$ identical controllers $\mathscr{C}_{i}=\mathscr{C}, i=$ $1,2, \ldots, N$, lead to consensus if and only if the controller $\mathscr{C}$ stabilizes the output variable $\bar{y}_{i}$ in all the plants $\mathscr{P}\left(\lambda_{i}\right) \wedge \mathscr{C}, i \in\{2, \ldots, N\}$, where $\lambda_{2}, \ldots, \lambda_{N}$ are the positive and real eigenvalues of $L$. Condition 4 . allows us to make use of some results previously derived for the dead-beat control of discrete-time systems [3], but whose adaption to the case of a stabilizing control (of part of the system variables) of a continuous-time system is straightforward.

Proposition 5. Given the behavior $\mathscr{P}\left(\lambda_{i}\right)$ described in (7), the controller $\mathscr{C}$ described in (4) stabilizes the output variable $\bar{y}_{i}$ in $\mathscr{P}\left(\lambda_{i}\right) \wedge \mathscr{C}$, (i.e., for every $\left(\frac{\bar{y}_{i}}{\bar{u}_{i}}\right) \in$ $\mathscr{P}\left(\lambda_{i}\right) \wedge \mathscr{C}$ the variable $\bar{y}_{i}$ is small) if and only if

$$
\Gamma_{i}(s):=\left(\begin{array}{ll}
A(s) & \lambda_{i} B(s) \tag{9}
\end{array}\right)\binom{P_{y}(s)}{C_{y}(s)}
$$

is of full column rank and the greatest common divisor of its maximal order minors


## 5 Consensus with regular interconnections

Up to now we have focused on the consensus target without imposing any constraint on the controller that allows to reach this goal. As a matter of fact, a natural constraint to introduce is that the connection of $\mathscr{\mathscr { P }}$ and $\mathscr{C}_{L}$ is well-posed, i.e., that the input/output structure of $\mathscr{P}$ is preserved even after interconnection with $\mathscr{C}_{L}$. This means that each agent may have the possibility to correct the signal obtained from the other agents via feedback interconnection, by adding some completely free component. As previously recalled, in Section 2, this goal is achieved by imposing that the interconnection of the agents and the controllers is regular. By the assumptions we have introduced on the descriptions of the agents and the controllers, the regularity of the interconnection $\mathscr{K}=\underline{\mathscr{P}} \wedge \underline{\mathscr{C}}_{L}$ amounts to requiring that the matrix

$$
K=\left(\begin{array}{ll}
I_{N} \otimes P_{y} & -I_{N} \otimes P_{u} \\
L \otimes C_{y} & -I_{N} \otimes C_{u}
\end{array}\right)
$$

is of full row rank, and $K$ being a square matrix, this means that it must be nonsingular. Equivalently, the connected system $\mathscr{K}$ is autonomous. The following proposition shows that regularity of this interconnection is equivalent to regularity of the other interconnections appearing in Theorem 4. In order to prove it, it is convenient to preliminarily notice that the matrix

$$
\begin{align*}
\bar{K} & =\left(\begin{array}{ll}
I_{N} \otimes P_{y} & -I_{N} \otimes P_{u} \\
D \otimes C_{y} & -I_{N} \otimes C_{u}
\end{array}\right) \\
& =\left(\begin{array}{cc}
V^{-1} \otimes I_{p} & 0 \\
0 & V^{-1} \otimes I_{m}
\end{array}\right) K\left(\begin{array}{cc}
V \otimes I_{p} & 0 \\
0 & V \otimes I_{m}
\end{array}\right) \tag{10}
\end{align*}
$$

where $V$ is the matrix of eigenvectors of $L$ given in (3), can be rewritten, by resorting to suitable row and column permutations, to the block diagonal form

$$
\bar{K}_{\text {block }}:=\left(\begin{array}{ccc}
K_{1} & &  \tag{11}\\
& \ddots & \\
& & K_{N}
\end{array}\right)
$$

where

$$
K_{i}:=\left(\begin{array}{cc}
P_{y} & -P_{u} \\
\lambda_{i} C_{y} & -C_{u}
\end{array}\right) \in \mathbb{R}[s]^{(p+m) \times(p+m)} .
$$

This immediately leads to the following result, whose proof is straightforward and hence omitted.

Proposition 6. The following conditions are equivalent:

1. The interconnection $\mathscr{K}=\underline{\mathscr{P}} \wedge \mathscr{C}_{L}$ is regular.
2. The interconnection $\mathscr{K}_{D}=\underline{\mathscr{P}} \wedge \underline{\mathscr{C}} D$ is regular
3. For every $i \in\{1, \ldots, N\}$, the interconnection $\mathscr{P} \wedge \mathscr{C}\left(\lambda_{i}\right)$ is regular.
4. For every $i \in\{1, \ldots, N\}$, the interconnection $\mathscr{P}\left(\lambda_{i}\right) \wedge \mathscr{C}$ is regular.

We can now restate the results of Theorem 4, by adding the requirement that the resulting interconnections are regular, i.e., that they are well-posed feedback interconnections. The proof is an immediate corollary of Theorem 4 and of Proposition 6 , and hence is omitted.

Theorem 7 (consensus with regularity). The following statements are equivalent:

1. (a) The interconnection $\mathscr{K}=\underline{\mathscr{P}} \wedge \underline{\mathscr{C}}_{L}$ is regular, and
(b) the controllers $\mathscr{C}_{i}=\mathscr{C}, i=1,2, \ldots, N$, lead to consensus;
2. (a) The interconnection $\mathscr{K}_{D}=\underline{\mathscr{P}} \wedge \underline{\mathscr{C}}_{D}$ is regular, and
(b) $\forall\binom{\bar{y}}{\bar{u}} \in \mathscr{K}_{D}, \forall i \in\{2, \ldots, N\}, \quad \bar{y}_{i} \quad$ is small.
3. For every $i \in\{2, \ldots, N\}$,
(a) the interconnections $\mathscr{P} \wedge \mathscr{C}\left(\lambda_{i}\right)$ are regular, and
(b) $\forall\binom{\bar{y}_{i}}{\bar{u}_{i}} \in \mathscr{P} \wedge \mathscr{C}\left(\lambda_{i}\right), \quad \bar{y}_{i} \quad$ is small.
4. For every $i \in\{2, \ldots, N\}$,
(a) the interconnections $\mathscr{P}\left(\lambda_{i}\right) \wedge \mathscr{C}$ are regular, and
(b) $\forall\binom{\bar{y}_{i}}{\bar{u}_{i}} \in \mathscr{P}\left(\lambda_{i}\right) \wedge \mathscr{C}, \quad \bar{y}_{i} \quad$ is small, i.e., $\mathscr{C}$ stabilizes the variable $\bar{y}_{i}$ in $\mathscr{P}\left(\lambda_{i}\right) \wedge \mathscr{C}$.

Clearly, condition 4. provides the characterization we want to explore in order to understand under what conditions on the original agents model the consensus problem is solvable. By making use of Proposition 5 and of the assumption on the matrices $P_{y}$ and $C_{u}$ involved in the agents and controller description, we can obtain algebraic characterizations of points (a) and (b) of condition 4. in Theorem 7. We will refer to these latter conditions as 4.(a) and 4.(b).

Proposition 8. Given the behavior $\mathscr{P}\left(\lambda_{i}\right)$ described as in (7), with $\operatorname{det} P_{y} \neq 0, a$ controller $\mathscr{C}$ described as in (4), with $\operatorname{det} C_{u} \neq 0$, is such that
4.(a) the interconnection $\mathscr{P}\left(\lambda_{i}\right) \wedge \mathscr{C}$ is regular, and
4.(b) $\forall\binom{\bar{y}_{i}}{\bar{u}_{i}} \in \mathscr{P}\left(\lambda_{i}\right) \wedge \mathscr{C}, \bar{y}_{i}$ is small, i.e. $\mathscr{C}$ stabilizes the variable $\bar{y}_{i}$ in $\mathscr{P}\left(\lambda_{i}\right) \wedge$ $\mathscr{C}$,
if and only if
a) $\operatorname{det}\left[C_{u}-\lambda_{i} C_{y} P_{y}^{-1} P_{u}\right] \neq 0$, and
b) $\Gamma_{i}=\left(\begin{array}{ll}A & \lambda_{i} B\end{array}\right)\binom{P_{y}}{C_{y}}=A P_{y}+\lambda_{i} B C_{y}$ is a square Hurwitz matrix, where $\left(\begin{array}{ll}A & B\end{array}\right)$ is an MLA of $\binom{P_{u}}{C_{u}}$.
Proof. We preliminarily notice that, by assumption, both $P_{y}$ and $C_{u}$ are square and nonsingular matrices, and the matrix $M_{i}$ takes the form:

$$
M_{i}=M_{i}(s):=\left(\begin{array}{cc}
P_{y}(s) & -\lambda_{i} P_{u}(s) \\
C_{y}(s) & -C_{u}(s)
\end{array}\right) \in \mathbb{R}[s]^{(m+p) \times(m+p)} .
$$

[Necessity] If the interconnection $\mathscr{P}\left(\lambda_{i}\right) \wedge \mathscr{C}$ is regular then

$$
\operatorname{rank}\left(M_{i}\right)=\operatorname{rank}\left(\begin{array}{ll}
P_{y} & -\lambda_{i} P_{u}
\end{array}\right)+\operatorname{rank}\left(\begin{array}{ll}
C_{y} & -C_{u}
\end{array}\right)=p+m
$$

namely $\operatorname{det} M_{i} \neq 0$. As $\operatorname{det} P_{y} \neq 0$, we can write

$$
\operatorname{det} M_{i}=\operatorname{det} P_{y} \cdot \operatorname{det}\left[C_{u}-\lambda_{i} C_{y} P_{y}^{-1} P_{u}\right] \neq 0
$$

and hence condition a) holds. On the other hand, if $\mathscr{C}$ stabilizes the variable $\bar{y}_{i}$ in $\mathscr{P}\left(\lambda_{i}\right) \wedge \mathscr{C}$, by Proposition 5 we can claim that $\Gamma_{i}$ is of full column rank and the greatest common divisor of its maximal order minors is Hurwitz. As we impose that $C_{u}$ is square and nonsingular, then by [4] the matrix $A$ must be square and nonsingular in turn, and hence $\Gamma_{i}$ is a $p \times p$ square matrix. So, condition b) holds.
[Sufficiency] If a) holds, then $\operatorname{det} M_{i} \neq 0$. This ensures that $M_{i}$ is square and nonsingular, and hence the the interconnection $\mathscr{P}\left(\lambda_{i}\right) \wedge \mathscr{C}$ is regular. If $b$ ) holds, then, in particular, $\Gamma_{i}$ is of full column rank and the greatest common divisor of its maximal order minors is Hurwitz. This ensures, by Proposition 5, that $\mathscr{C}$ stabilizes the variable $\bar{y}_{i}$ in $\mathscr{P}\left(\boldsymbol{\lambda}_{i}\right) \wedge \mathscr{C}$.

We want to prove that, as a consequence of the previous proposition, if the consensus problem is solvable, then the output $y_{i}$ is stabilizable from the input $u_{i}$ in $\mathscr{P}_{i}=\mathscr{P}$.

Corollary 9. Suppose that the $N$ agents are described by the model (2), with $P_{y}(s) \in \mathbb{R}[s]^{p \times p}$ and $P_{u}(s) \in \mathbb{R}[s]^{p \times m}$, $\operatorname{det} P_{y} \neq 0$, and let $0<\lambda_{2} \leq \lambda_{3} \leq \cdots \leq \lambda_{N}$ be the positive real eigenvalues of $L$. If the consensus problem is solvable, namely polynomial matrices $C_{y}(s)$ and $C_{u}(s)$, of suitable sizes, with $\operatorname{det} C_{u} \neq 0$, can be found such that conditions $a$ ) and b) of Proposition 8 hold, then $y_{i}$ is stabilizable from $u_{i}$ in $\mathscr{P}_{i}=\mathscr{P}$.

Proof. If $P_{u}$ is of full row rank the result follows immediately from Proposition 1.
 following way:

$$
(A(s) \quad B(s))=\left(\begin{array}{cc}
M_{u}(s) & 0 \\
Y_{p}(s) & Y_{c}(s)
\end{array}\right)
$$

where $M_{u} \in \mathbb{R}[s]^{(p-r) \times p}$ is an MLA of $P_{u}$, and $Y_{p}$ and $Y_{c}$ are suitably chosen polynomial matrices. So, w.l.o.g. we can assume that

$$
\Gamma_{i}(s)=A(s) P_{y}(s)+\lambda_{i} B(s) C_{y}(s)=\binom{M_{u}(s) P_{y}(s)}{Y_{p}(s) P_{y}(s)+\lambda_{i} Y_{c}(s) C_{y}(s)}
$$

Being of full row rank, the $(p-r) \times p$ matrix $M_{u}(s) P_{y}(s)$ factorizes as

$$
M_{u}(s) P_{y}(s)=\Delta(s) R_{y}(s)
$$

where $\Delta$ is square and nonsingular and $R_{y}$ is left prime. So, if b) holds, and hence $\Gamma_{i}$ is a square Hurwitz matrix, then $\Delta$ must be Hurwitz in turn, thus proving that $y_{i}$ in stabilizable from $u_{i}$ in $\mathscr{P}_{i}=\mathscr{P}$ (see Proposition 1).

## 6 The problem solution

In the previous section we have shown that the consensus problem among $N$ identical agents, each of them described as in (2), with the regularity constraint on the agents/controllers interconnection, is equivalent to the problem of determining polynomial matrices $C_{y}(s) \in \mathbb{R}[s]^{m \times p}$ and $C_{u}(s) \in \mathbb{R}[s]^{m \times m}$, with $\operatorname{det} C_{u} \neq 0$, such that conditions $a$ ) and $b$ ) of Proposition 8 hold. Also, if this is the case, then for each $i$ th agent $\mathscr{P}_{i}=\mathscr{P}$ the output $y_{i}$ is stabilizable from the input $u_{i}$.

In this section we want to prove that this necessary assumption on the agents $\mathscr{P}_{i}, i=1, \ldots, N$, is also sufficient for the consensus problem to be solvable. To prove this, we need a preliminary lemma.

Lemma 10. Given a nonzero polynomial $\gamma(s) \in \mathbb{R}[s]$ and positive real numbers $0<\lambda_{2} \leq \cdots \leq \lambda_{N}$, a polynomial $n(s) \in \mathbb{R}[s]$ can be found such that the $N-1$ polynomials $\gamma(s)+\lambda_{i} n(s), i=2, \ldots, N$, are Hurwitz.

Proof. Factorize $\gamma(s)$ as $\gamma(s)=\gamma_{s}(s) \gamma_{u}(s)$, where $\gamma_{s}(s)$ is Hurwitz, while $\gamma_{u}(s)$ is monic and has all its roots in the closed right half-plane $\{s \in \mathbb{C}: \operatorname{Re}(s) \geq 0\}$. Also, set $d:=\operatorname{deg} \gamma_{u}(s)$. It is easily seen, by resorting to the positive root locus plot, that there exists $\bar{K}>0$ such that for every $K \geq \bar{K}$, the polynomial $\gamma_{u}(s)+K(s+1)^{d}$ is Hurwitz. This ensures that by assuming

$$
n(s)=\frac{\bar{K}}{\lambda_{2}} \gamma_{s}(s)(s+1)^{d}
$$

we have that all polynomials $\gamma(s)+\lambda_{i} n(s), i=2, \ldots, N$, are Hurwitz.

In the following, we address two cases: 1) the case when $P_{u}$ is of full row rank; 2) the case when rank $P_{u}=: r<p$.

### 6.1 Case 1: $P_{u}$ is of full row rank

If $P_{u}$ is of full row rank, then (see Proposition 1) $y_{i}$ is necessarily stabilizable from $u_{i}$. Let $P_{u c}$ be a completion of $P_{u}$ to a square and nonsingular matrix, and set

$$
C_{u}:=\binom{P_{u}}{P_{u c}} \in \mathbb{R}[s]^{m \times m} .
$$

Clearly, $\operatorname{det} C_{u} \neq 0$. Moreover, it is immediately seen that

$$
\binom{P_{u}}{C_{u}}=\left(\begin{array}{cc}
I_{p} & 0 \\
I_{p} & 0 \\
0 & I_{m-p}
\end{array}\right) C_{u},
$$

and hence an MLA of $\binom{P_{u}}{C_{u}}$ is

$$
\left(\begin{array}{ll}
A & B
\end{array}\right)=\left(\begin{array}{ll}
I_{p} & -I_{p} \\
0_{p \times(m-p)}
\end{array}\right) .
$$

This implies that

$$
A P_{y}+\lambda_{i} B C_{y}=P_{y}-\lambda_{i} C_{y A},
$$

where $C_{y A} \in \mathbb{R}[s]^{p \times p}$ is the block of the first $p$ rows of $C_{y}$. Let $U \in \mathbb{R}[s]^{p \times p}$ be a unimodular matrix such that $U P_{y}$ is in lower triangular Hermite form [9], and let $\gamma_{1}(s), \ldots, \gamma_{p}(s)$ be its diagonal entries. By Lemma 10 , polynomials $n_{1}(s), \ldots, n_{p}(s)$ can be found such that the $p(N-1)$ polynomials $\gamma_{j}(s)+\lambda_{i} n_{j}(s), j=1, \ldots, p, i=$ $2, \ldots, N$, are Hurwitz. Consequently, by assuming

$$
C_{y A}=-U^{-1} \operatorname{diag}\left\{n_{1}, \ldots, n_{p}\right\}
$$

(and by arbitrarily choosing the remaining $m-p$ rows of $C_{y}$ ) we ensure that condition b) of Proposition 8 holds. It is trivial to verify that also condition a) holds.

### 6.2 Case 2: $P_{u}$ is of rank $r<p$

Assume that rank $P_{u}=: r<p$ and factorize $P_{u}$ as $P_{u}=L_{u} R_{u}$, where $L_{u} \in \mathbb{R}[s]^{p \times r}$ is right prime and $R_{u} \in \mathbb{R}[s]^{r \times m}$ is of full row rank. Consider a completion $R_{u c}$ of $R_{u}$ to a nonsingular $m \times m$ matrix ${ }^{1}$, and set

$$
C_{u}:=\binom{R_{u}}{R_{u c}} \in \mathbb{R}[s]^{m \times m} .
$$

Clearly, $\operatorname{det} C_{u} \neq 0$. Moreover, it is immediately seen that

$$
\binom{P_{u}}{C_{u}}=\left(\begin{array}{cc}
L_{u} & 0 \\
I_{r} & 0 \\
0 & I_{m-r}
\end{array}\right) C_{u},
$$

[^0]and hence an MLA of $\binom{P_{u}}{C_{u}}$ is
\[

\left($$
\begin{array}{ll}
A & B
\end{array}
$$\right)=\left($$
\begin{array}{lll}
I_{p} & -L_{u} & 0_{p \times(m-r)}
\end{array}
$$\right) .
\]

This implies that

$$
A P_{y}+\lambda_{i} B C_{y}=P_{y}-\lambda_{i} L_{u} C_{y A}
$$

where $C_{y A}$ is the block of the first $r$ rows of $C_{y}$. Let

$$
U=\binom{U_{1}}{U_{2}} \begin{aligned}
& \} r \\
& \} p-r
\end{aligned} \in \mathbb{R}[s]^{p \times p}
$$

be a unimodular matrix such that

$$
U L_{u}=\binom{I_{r}}{0_{(p-r) \times r}} .
$$

Note that $U_{2}$ is an MLA of $L_{u}$ and hence of $P_{u}$. On the other hand, set

$$
\binom{U_{1}}{U_{2}} P_{y}=:\binom{U_{1} P_{y}}{\Delta_{2} R_{y}}
$$

The assumption that $y_{i}$ is stabilizable from $u_{i}$ in $\mathscr{P}_{i}=\mathscr{P}$, together with the fact that $U_{2} P_{y}$ is of full row rank, ensure that $U_{2} P_{y}$ factorizes as $U_{2} P_{y}=\Delta_{2} R_{y}$, where $\Delta_{2}$ is square and nonsingular and Hurwitz, while $R_{y}$ is left prime (see Proposition 1). This ensures that

$$
\begin{aligned}
U\left(A P_{y}+\lambda_{i} B C_{y}\right) & =U\left(P_{y}-\lambda_{i} L_{u} C_{y A}\right)=\left(\begin{array}{cc}
I_{r} & 0 \\
0 & \Delta_{2}
\end{array}\right)\binom{U_{1} P_{y}}{R_{y}}-\lambda_{i}\binom{C_{y A}}{0} \\
& =\left(\begin{array}{cc}
I_{r} & 0 \\
0 & \Delta_{2}
\end{array}\right)\binom{U_{1} P_{y}-\lambda_{i} C_{y A}}{R_{y}}
\end{aligned}
$$

We want to show that it is possible to choose $C_{y A}$ in such a way that the matrix $U\left(P_{y}-\lambda_{i} L_{u} C_{y A}\right)$, and hence $A P_{y}+\lambda_{i} B C_{y}$, will necessarily be square and Hurwitz for every $i=2, \ldots, N$. Let $V$ be a $p \times p$ unimodular matrix such that

$$
\binom{U_{1} P_{y}}{R_{y}} V=\left(\begin{array}{cc}
X_{1} & X_{2} \\
0 & I_{p-r}
\end{array}\right)
$$

and $X_{1} \in \mathbb{R}[s]^{r \times r}$ is in lower triangular Hermite form, with $\gamma_{1}(s), \ldots, \gamma_{r}(s)$ as diagonal entries. By Lemma 10, polynomials $n_{1}(s), \ldots, n_{r}(s)$ can be found such that the $p(N-1)$ polynomials $\gamma_{j}(s)+\lambda_{i} n_{j}(s), j=1, \ldots, r, i=2, \ldots, N$, are Hurwitz. Consequently, by assuming

$$
C_{y A}=-\left(\operatorname{diag}\left\{n_{1}, \ldots, n_{r}\right\} \quad 0_{r \times(p-r)}\right) V^{-1}
$$

(while the lower block of rows in $C_{y}$ is arbitrary), we ensure that condition b) of Proposition 8 holds. It is trivial to verify that also condition a) holds.

So, to summarize, by putting together the results of this section with Corollary 9 , we obtain the following result.

Theorem 11. Suppose that the $N$ agents are described by the model (2), with $P_{y}(s) \in \mathbb{R}[s]^{p \times p}$ and $P_{u}(s) \in \mathbb{R}[s]^{p \times m}$, $\operatorname{det} P_{y} \neq 0$, and let $0<\lambda_{2} \leq \lambda_{3} \leq \cdots \leq \lambda_{N}$ be the positive real eigenvalues of $L$. The following facts are equivalent:
i) $y_{i}$ in stabilizable from $u_{i}$ in $\mathscr{P}_{i}=\mathscr{P}$,
ii) the consensus problem is solvable, namely polynomial matrices $C_{y}(s)$ and $C_{u}(s)$, of suitable sizes, with $\operatorname{det} C_{u} \neq 0$, can be found such that conditions a) and b) of Proposition 8 hold.

## 7 Comparison with the state-space set-up

In this section we aim at comparing the set-up and the results derived in the previous sections with the most common set-up and related results available in the literature. We refer for the comparison to [20], but several other references could be mentioned (see, e.g., [7, 14, 21] and references therein). In [20] each agent is described by a state-space model and hence the dynamics takes the form

$$
\begin{aligned}
\mathscr{P}_{i}:=\mathscr{P}: & =\left\{\binom{x_{i}}{u_{i}} \in \mathscr{F}^{n+m}: \dot{x}_{i}=A x_{i}+B u_{i}\right\} \\
& =\left\{\binom{x_{i}}{u_{i}} \in \mathscr{F}^{n+m}:\left(s I_{n}-A\right) \circ x_{i}=B \circ u_{i}\right\}
\end{aligned}
$$

for constant matrices $A \in \mathbb{R}^{n \times n}$ and $B \in \mathbb{R}^{n \times m}$. This amounts to assuming

$$
\left(\begin{array}{ll}
P_{y} & -P_{u}
\end{array}\right)=\left(\begin{array}{ll}
s I-A & -B
\end{array}\right)
$$

and clearly $\operatorname{det} P_{y}=\operatorname{det}(s I-A) \neq 0$. The agents' network is described by a directed weighted graph. The information available to the $i$ th agent is

$$
\widetilde{x}_{i}:=\sum_{j=1}^{N} a_{i j}\left(x_{j}-x_{i}\right)=-\sum_{j=1}^{N} a_{i j}\left(x_{i}-x_{j}\right)
$$

(which differs from the present set-up just in the sign). The consideblack controllers are of the form

$$
\begin{aligned}
\mathscr{C}_{i}:=\mathscr{C} & :=\left\{\binom{\widetilde{x}_{i}}{u_{i}} \in \mathscr{F}^{n+m}: u_{i}=K \widetilde{x}_{i}\right\} \\
& =\left\{\binom{\widetilde{x}_{i}}{u_{i}} \in \mathscr{F}^{n+m}: K \circ \widetilde{x}_{i}=I_{m} \circ u_{i}\right\}
\end{aligned}
$$

for a constant matrix $K \in \mathbb{R}^{m \times n}$. Note that, by the structure imposed on the feedback control law in [20], $C_{u}=I_{m}$, while $C_{y}$ is replaced by a constant matrix. This automatically ensures that

$$
\operatorname{rank}\left(\begin{array}{cc}
P_{y} & -\lambda_{i} P_{u} \\
C_{y} & -C_{u}
\end{array}\right)=\operatorname{rank}\left(\begin{array}{ll}
P_{y} & -\lambda_{i} P_{u}
\end{array}\right)+\operatorname{rank}\left(\begin{array}{ll}
C_{y} & -C_{u}
\end{array}\right),
$$

namely the interconnection is regular, for every value of $\lambda_{i}$.
Lemma 3.1 in [20], which is quoted from [7, 21], states that consensus on $x$ is reached if and only if the matrices $A-\lambda_{i} B K, i=2, \ldots, N$, are Hurwitz. This is equivalent to the condition

$$
\forall i \in\{2, \ldots, N\}, \forall\binom{\bar{x}_{i}}{\bar{u}_{i}} \in \mathscr{P} \wedge \mathscr{C}\left(-\lambda_{i}\right), \quad \bar{x}_{i} \text { is small, }
$$

since

$$
\begin{aligned}
\mathscr{P} \wedge \mathscr{C}\left(-\lambda_{i}\right) & =\left\{\binom{\bar{x}_{i}}{\bar{u}_{i}} \in \mathscr{F}^{n+m}:\left(s I_{n}-A\right) \circ \bar{x}_{i}=B \circ \bar{u}_{i}, \bar{u}_{i}=-\lambda_{i} K \circ \bar{x}_{i}\right\} \\
& =\left\{\binom{\bar{x}_{i}}{\bar{u}_{i}} \in \mathscr{F}^{n+m}:\left(s I_{n}-A\right) \circ \bar{x}_{i}=-\lambda_{i} B K \circ \bar{x}_{i}, \bar{u}_{i}=-\lambda_{i} K \circ \bar{x}_{i}\right\} \\
& =\left\{\binom{\bar{x}_{i}}{\bar{u}_{i}} \in \mathscr{F}^{n+m}:\left(s I_{n}-\left(A-\lambda_{i} B K\right)\right) \circ \bar{x}_{i}=0, \bar{u}_{i}=-\lambda_{i} K \circ \bar{x}_{i}\right\} .
\end{aligned}
$$

Taking into account the different definition of $\widetilde{x}_{i}$ leading to the minus sign in $\mathscr{C}\left(-\lambda_{i}\right)$, and the fact that in [20] the target variables are $x_{i}$ while here we refer to $y_{i}$, this result is in perfect agreement with Theorem 4. In addition, in Theorem 3.3 of [20] it is stated that, under the assumption that the communication graph is connected, the consensus problem is solvable if and only if the agent's state-space model is stabilizable. So, the theory derived in this paper in the behavioral context naturally extends and generalizes the state-space approach previously investigated in the literature.

It is worth mentioning that in general the fact that the $N$ agents reach consensus, and hence that for every $\binom{\bar{y}_{i}}{\bar{u}_{i}} \in \mathscr{P}\left(\lambda_{i}\right) \wedge \mathscr{C}, i=2, \ldots, N$, the variable $\bar{y}_{i}$ is small, does not impose any constraint on the corresponding input $\bar{u}_{i}$ that, on the contrary, may diverge as $t \rightarrow+\infty$. Clearly, this situation never arises when dealing with state-space models since one always resorts to a static state-feedback, and hence the convergence to zero of the state trajectory naturally ensures that the corresponding input converges to zero, in turn. If we want to reach consensus also on the input trajectories, then we need to impose the stabilizability property on both variables in $\mathscr{P}_{i}=\mathscr{P}$.

## References

[1] C. Altafini. Consensus problems on networks with antagonistic interactions. IEEE Trans. Automat. Control, 58 (4):935-946, 2013.
[2] M. N. Belur and H. L. Trentelman. Stabilization, pole placement, and regular implementability. IEEE Trans. Automat. Control, 47 (5):735-744, 2002.
[3] M. Bisiacco and M. E. Valcher. Dead-beat control in the behavioral approach. IEEE Trans. Automat. Control, 57 (9):2168-2175, 2012.
[4] M. Bisiacco and M.E. Valcher. Behavior decompositions and two-sided diophantine equations. Automatica, 37 (9):1387-1395, 2001.
[5] I. Blumthaler, M. Bisiacco and M.E. Valcher. The consensus problem in the behavioral approach. Proc. 53rd IEEE Conf. Decision and Control, pages 727-734,Los Angeles (CA), 2014.
[6] M.H. DeGroot. Reaching a consensus. J. Amer. Statist. Assoc., 69 (345):118121, 1974.
[7] J. A. Fax and R. M. Murray. Information flow and cooperative control of vehicle formations. IEEE Trans. Automat. Control, 49 (9):1465-1476, 2004.
[8] A. Jadbabaie, J. Lin, and A.S. Morse. Coordination of groups of mobile autonomous agents using nearest neighbor rules. IEEE Trans. Automat. Control, 48 (6):988-1001, 2003.
[9] T. Kailath. Linear Systems. Prentice-Hall, 1980.
[10] J. Lin, A.S. Morse, and B.D.O. Anderson. The multi-agent rendezvous problem. In Proc. 42nd IEEE Conf. Decision and Control, pages 1508 -1513, Maui, Hawaii, 2003.
[11] L. Moreau. Stability of multiagent systems with time-dependent communication links. IEEE Trans. Automat. Control, 50 (2):169-182, 2005.
[12] R. Olfati-Saber, J. A. Fax, and R. M. Murray. Consensus and cooperation in networked multi-agent systems. Proc. IEEE, 95 (1):215-223, 2007.
[13] J. W. Polderman and J. C. Willems. Introduction to Mathematical Systems Theory. A Behavioral Approach. Springer-Verlag, New York, 1998.
[14] W. Ren and R. W. Beard. Consensus seeking in multiagent systems under dynamically changing interaction topologies. IEEE Trans. Automat. Control, 50 (5):655-661, 2005.
[15] P. Rocha. Structure and Representation of 2-D Systems. PhD thesis, University of Groningen, The Netherlands, 1990.
[16] L. Scardovi and R. Sepulchre. Synchronization in networks of identical linear systems. Automatica, 45 (11):2557-2562, 2009.
[17] H.L. Smith. Monotone Dynamical Systems: An Introduction to the Theory of Competitive and Cooperative Systems. AMS, Mathematical Surveys and Monographs, Providence, RI, vol. 41, 1995.
[18] J.N. Tsitsiklis. Problems in Decentralized Decision Making and Computation. PhD thesis, Department of EECS, MIT, 1984.
[19] S.E. Tuna. LQR-based coupling gain for synchronization of linear systems, arXiv:0801.3390v1 [math.OC], 2008.
[20] P. Wieland, J.-S. Kim, and F. Allgöwer. On topology and dynamics of consensus among linear high-order agents. Internat. J. Systems Sci., 42 (10):18311842, 2011.
[21] P. Wieland, J.-S. Kim, H. Scheu, and F. Allgöwer. On consensus in multi-agent systems with linear high-order agents. Proc. 17th IFAC World Congress, pages 1541-1546, 2008.
[22] J. C. Willems. From time series to linear system. I. Finite-dimensional linear time invariant systems. Automatica, 22 (5):561-580, 1986.
[23] J. C. Willems. Paradigms and puzzles in the theory of dynamical systems. IEEE Trans. Automat. Control, 36 (3):259-294, 1991.
[24] J. C. Willems. On interconnections, control, and feedback. IEEE Trans. Automat. Control, 42 (3):326-339, 1997.


[^0]:    ${ }^{1}$ If $r=m$ then $R_{u}$ is already square and hence we assume $C_{u}=R_{u}$. All the subsequent reasonings remain valid.

