



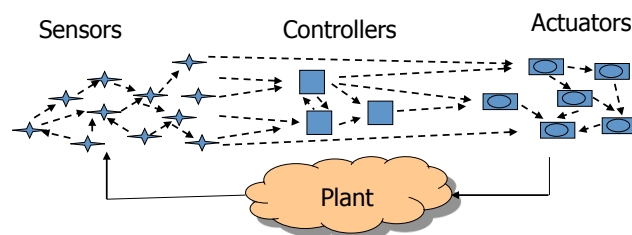
Networked Control and Autonomy

Karl Henrik Johansson
KTH Royal Institute of Technology
Stockholm, Sweden

2013 Summer School of Information Engineering
June 30 – July 6, 2013 in Bressanone, Italy

Slides available at <http://www.ee.kth.se/~kallej>

Networked control system





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Outline

Lecture 1: Motivating applications and challenges

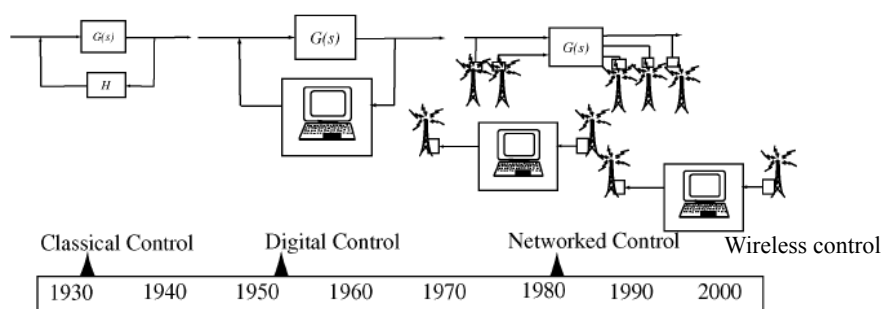
Lecture 2: Wireless control systems



Lecture 2 Outline

- What's new with wireless networked control?
- State-based scheduling for control
- Exploiting wireless protocols for control
- Event-based control
- Conclusions

A history of control

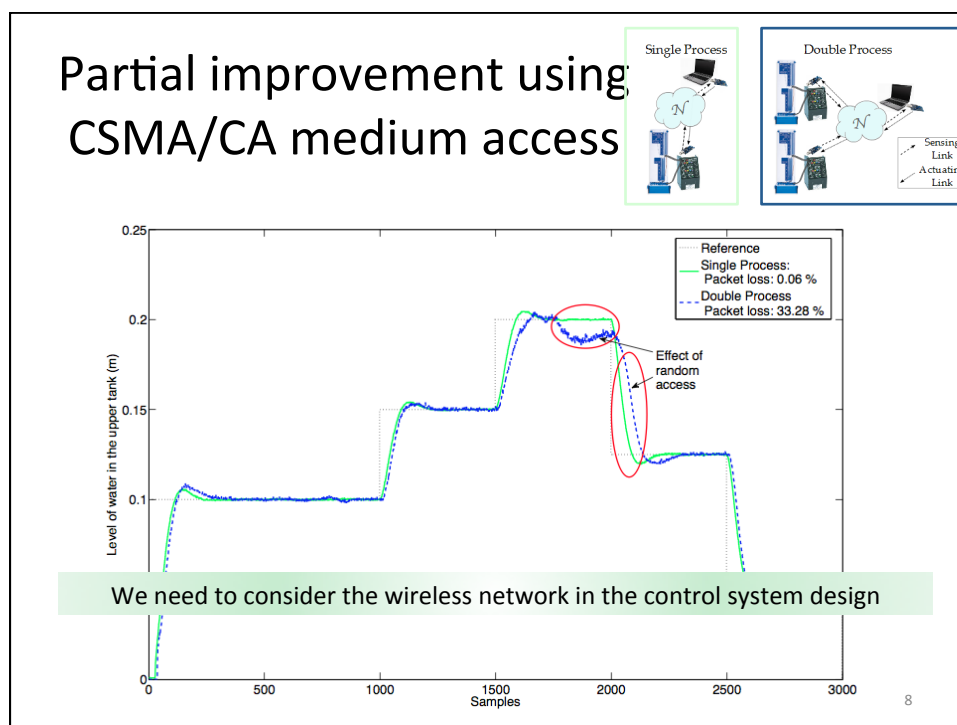
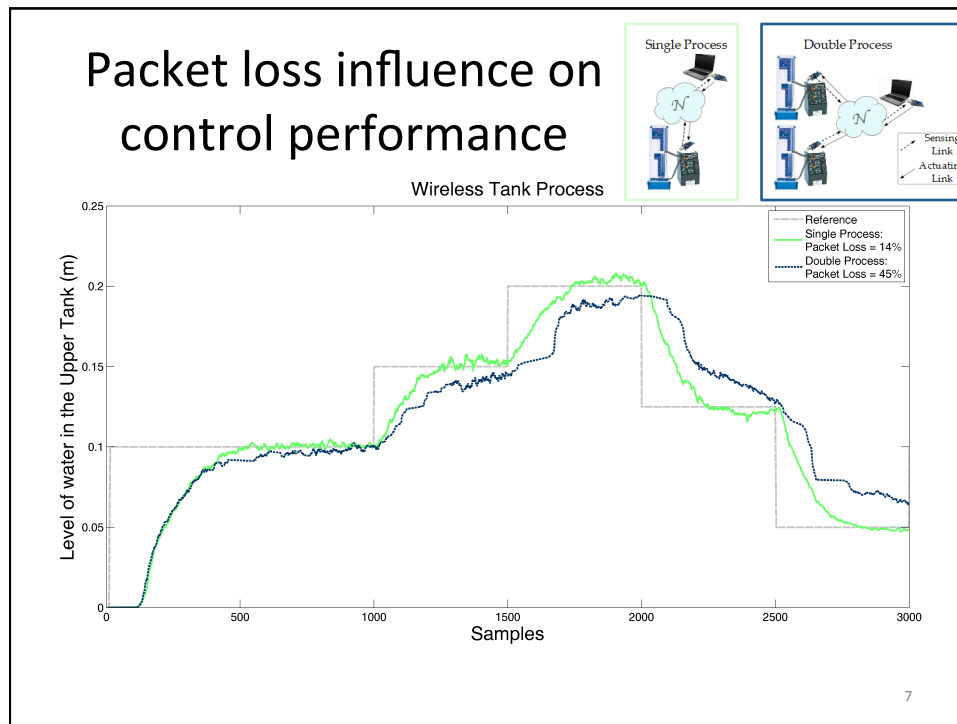


From dedicated communication links and networks for control systems
To open and ubiquitous wireless networks for control applications

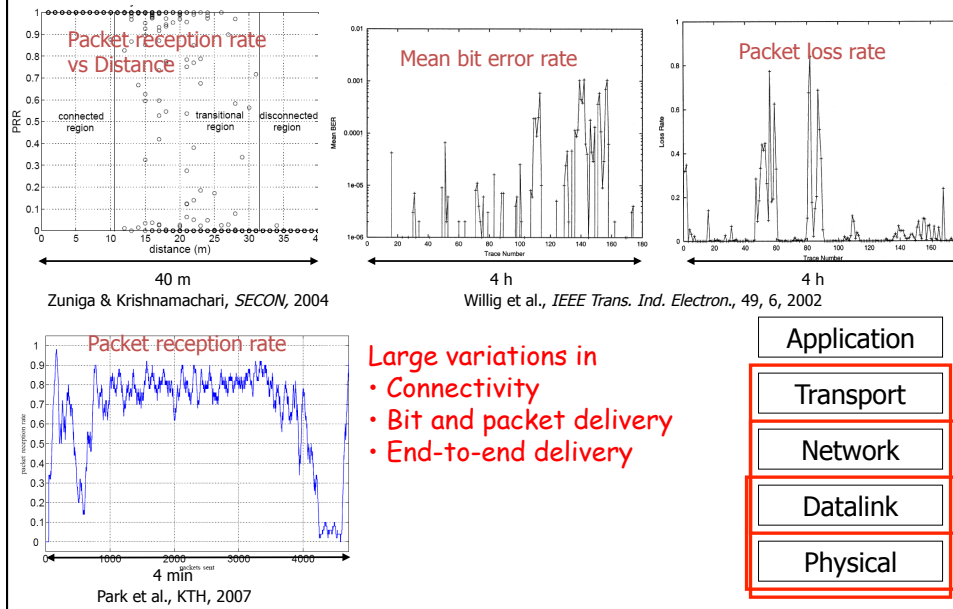
Adopted from [Baillieul & Antsaklis, 2007]

How share common network resources while maintaining guaranteed closed-loop performance?



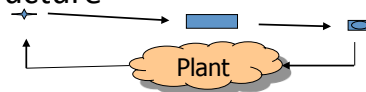


Uncertainty on several communication layers



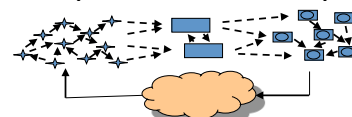
What's new with control over wireless networks?

- **Traditional control** systems design is based on assumption of **perfect information** being circulated in the system
- Information flow are dedicated to specific control loop
- Devices need to be fixed to infrastructure



Wireless control systems have

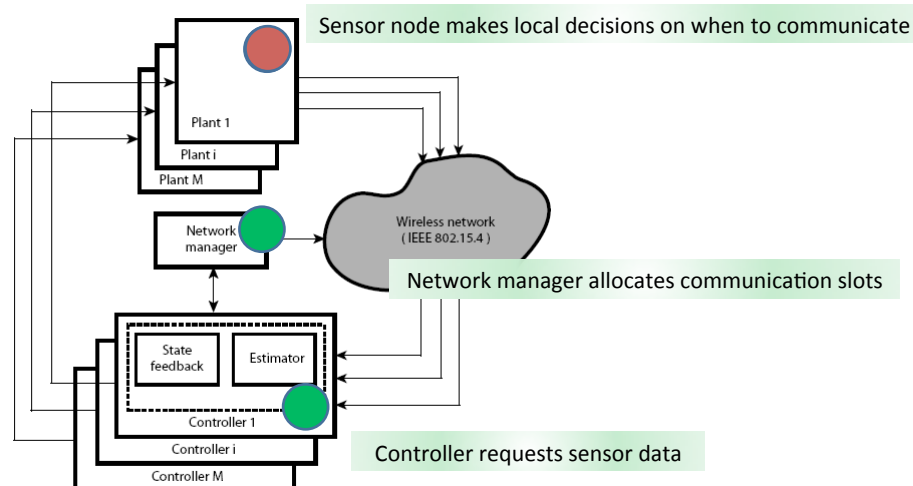
- **Non-ideal communication** between system devices; leads to interference, congestion, delay, loss, outages etc.
- + **Information** can be **shared** between components and loops
- + Enhanced **mobility and flexibility**



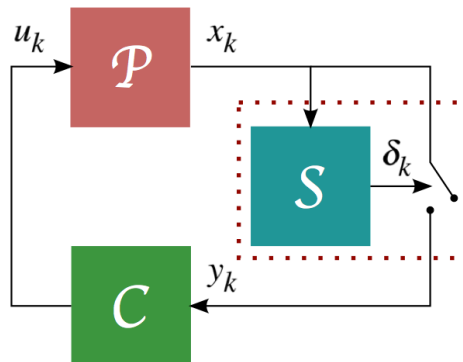
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Where taking medium access decisions?



Is there a separation between scheduling-estimation-control?



Stochastic control formulation

Plant:

$$x_{k+1} = Ax_k + Bu_k + w_k$$

Scheduler:

$$\delta_k = f_k(\mathbb{I}_k^S) \in \{0, 1\}$$

$$\mathbb{I}_k^S = [\{x\}_0^k, \{y\}_0^{k-1}, \{\delta\}_0^{k-1}, \{u\}_0^{k-1}]$$

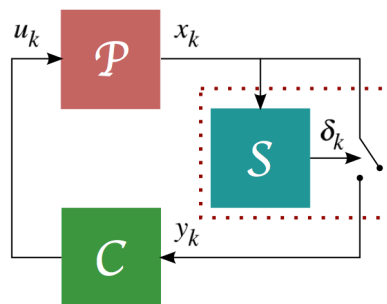
Controller:

$$u_k = g_k(\mathbb{I}_k^C)$$

$$\mathbb{I}_k^C = [\{y\}_0^k, \{\delta\}_0^k, \{u\}_0^{k-1}]$$

Cost criterion:

$$J(f, g) = \mathbb{E}[x_N^T Q_0 x_N + \sum_{s=0}^{N-1} (x_s^T Q_1 x_s + u_s^T Q_2 u_s)]$$



Control without scheduling = Classical LQG control of Kalman

The controller minimizing

$$J = \mathbb{E} \left[x_N^T Q_0 x_N + \sum_{s=0}^{N-1} (x_s^T Q_1 x_s + u_s^T Q_2 u_s) \right]$$

is given by

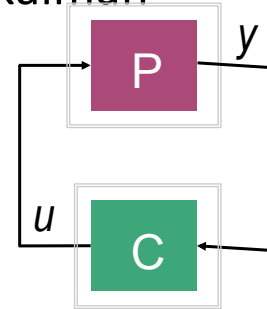
$$u_k = -L_k \hat{x}_{k|k},$$

$$L_k = (Q_2 + B^T S_{k+1} B)^{-1} B^T S_{k+1} A$$

where

$$S_k = Q_1 + A^T S_{k+1} A - A^T S_{k+1} B (Q_2 + B^T S_{k+1} B)^{-1} B^T S_{k+1} A$$

$\hat{x}_{k|k} = \mathbb{E}[x_k | \{y\}_0^k, \{u\}_0^{k-1}]$ is the minimum mean-square error (MMSE) estimate



Kalman, 1960

Event-based scheduler

Plant:

$$x_{k+1} = Ax_k + Bu_k + w_k$$

Scheduler:

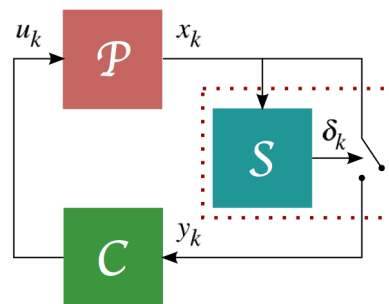
$$\delta_k = f_k(\mathbb{I}_k^S) \in \{0, 1\}$$

$$\mathbb{I}_k^S = [\{x\}_0^k, \{y\}_0^{k-1}, \{\delta\}_0^{k-1}, \{u\}_0^{k-1}]$$

Controller:

$$u_k = g_k(\mathbb{I}_k^C)$$

$$\mathbb{I}_k^C = [\{y\}_0^k, \{\delta\}_0^k, \{u\}_0^{k-1}]$$



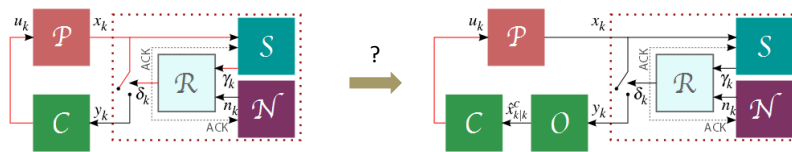
The separation principle does not hold for the optimal closed-loop system, so the design of the (event-based) scheduler, estimator, and controller is coupled

Feldbaum, 1965; Åström, 1970; Bar-Shalom and Tse, 1974

Ramesh et al., 2011

Conditions for Separation

Corollary: The optimal controller for the system $\{\mathcal{P}, S(f), C(g)\}$, with respect to the cost J is certainty equivalent if and only if the scheduling decisions are not a function of the applied controls.



Nice architecture achieved at the cost of optimality

Ramesh et al., 2011

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Event-based control architecture

- Plant \mathcal{P} :

$$x_{k+1} = ax_k + bu_k + w_k$$

- State-based Scheduler \mathcal{S} :

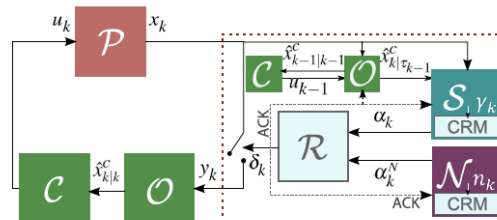
$$\gamma_k = \begin{cases} 1, & |x_k - \hat{x}_{k|\tau_{k-1}}^s|^2 > \epsilon_d, \\ 0, & \text{otherwise.} \end{cases}$$

$$\hat{x}_{k|\tau_{k-1}}^s = a\hat{x}_{k-1|k-1}^c + bu_{k-1}$$

- CRM: $\mathbb{P}(\alpha_k=1|\gamma_k=1) = \mathbb{P}(\alpha_k^N=1|n_k=1) = p_\alpha$
 $\delta_k = \alpha_k(1 - \alpha_k^N)$

- Observer \mathcal{O} : $y_k^{(j)} = \delta_k^{(j)} x_k^{(j)}$
 $\hat{x}_{k|k}^c = \bar{\delta}_k(a\hat{x}_{k-1|k-1}^c + bu_{k-1}) + \delta_k x_k$

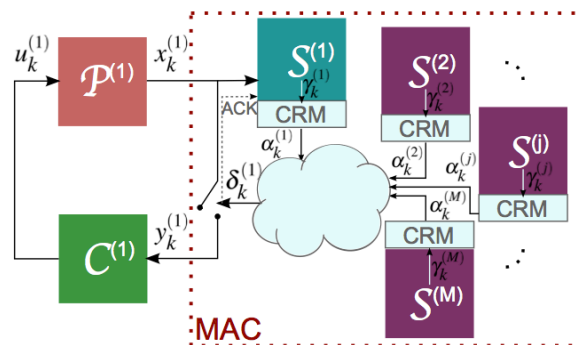
- Controller \mathcal{C} : $u_k = -L\hat{x}_{k|k}^c$



Ramesh et al., CDC, 2012

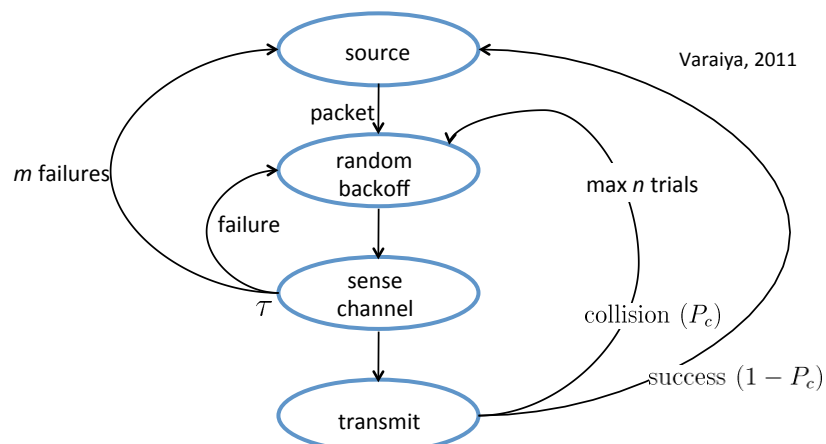
How to integrate contention resolution mechanisms?

- Hard problem because of correlation between transmissions (and the plant states)
- Closed-loop analysis can still be done for classes of event-based schedulers and MAC's



Ramesh et al., CDC 2011

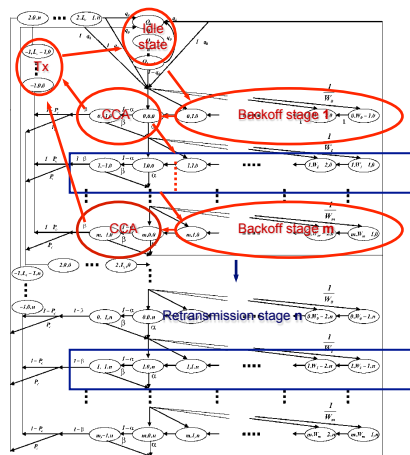
Contention resolution through CSMA/CA



- Every transmitting device executes this protocol
- For analysis, assume carrier sense events are independent [Bianchi, 2000]

CSMA/CA = Carrier Sense Multiple Access with Collision Avoidance

Detailed model of CSMA/CA in IEEE 802.15.4



- Markov state (s, c, r)
 - s : backoff stage
 - c : state of backoff counter
 - r : state of retransmission counter
 - Model parameters
 - q_0 : traffic condition ($q_0=0$ saturated)
 - m_0, m, m_b, n : MAC parameters
 - Computed characteristics
 - α : busy channel probability during CCA1
 - β : busy channel probability during CCA2
 - P_c : collision probability
- Detailed model for numerical evaluations
 - Reduced-order models for control design
 - Validated in simulation and experiment

Park, Di Marco, Soldati, Fischione, J, 2009

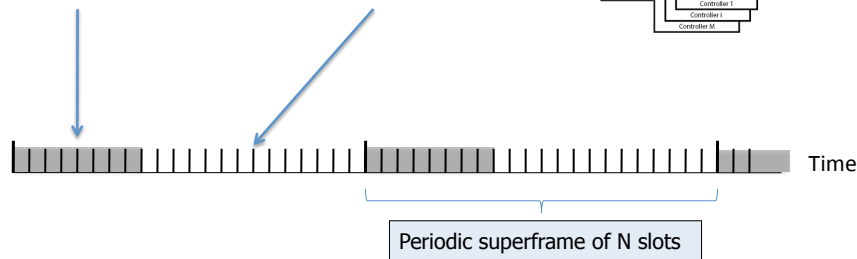
Cf., Bianchi, 2000; Pollin et al., 2006

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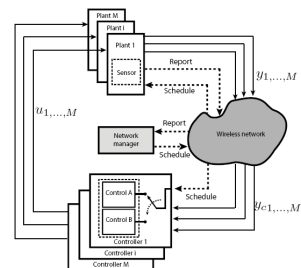
Slotted medium access

Many medium access protocols have slotted contention-free and contention access periods

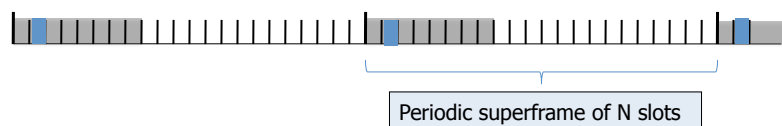


Hybrid MAC protocols

Exploit the mix of CFP's and CAP's for networked control



Contention-free period for TDMA scheduled communication

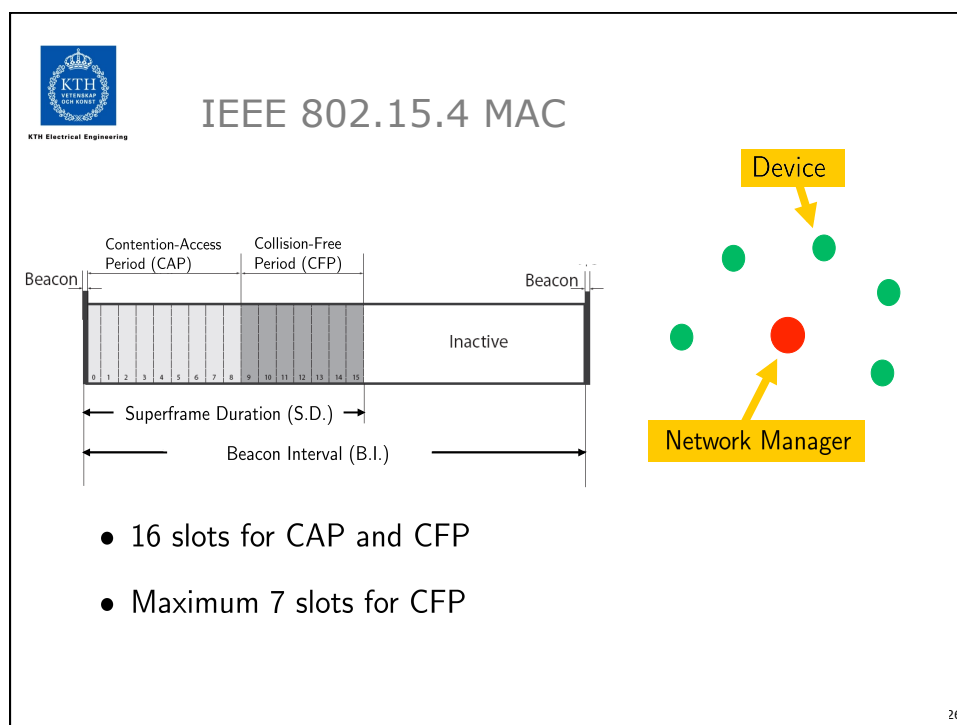
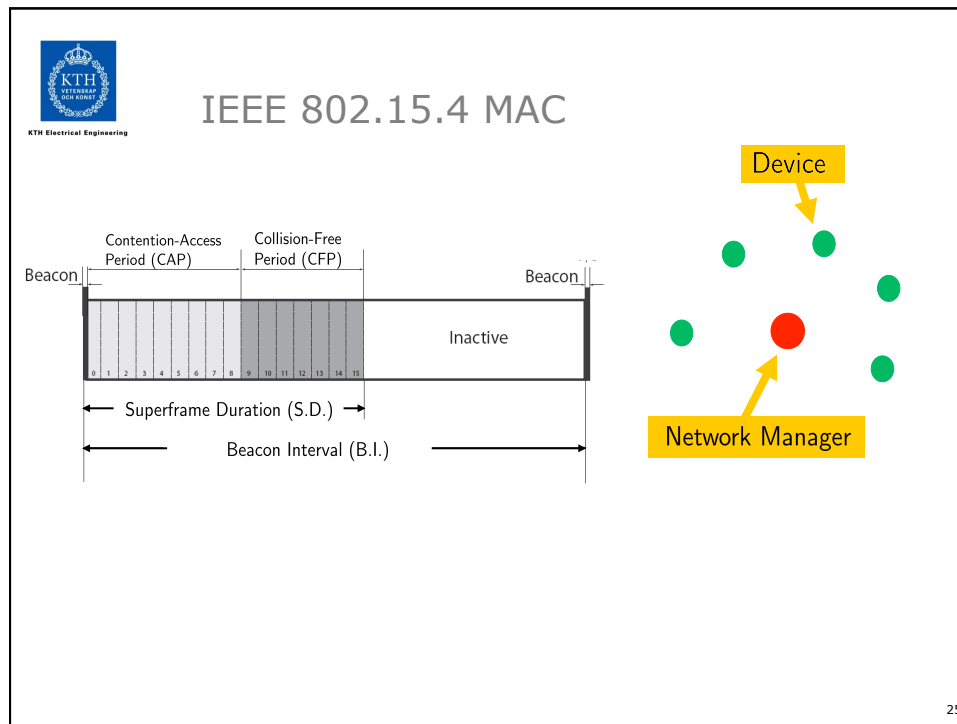


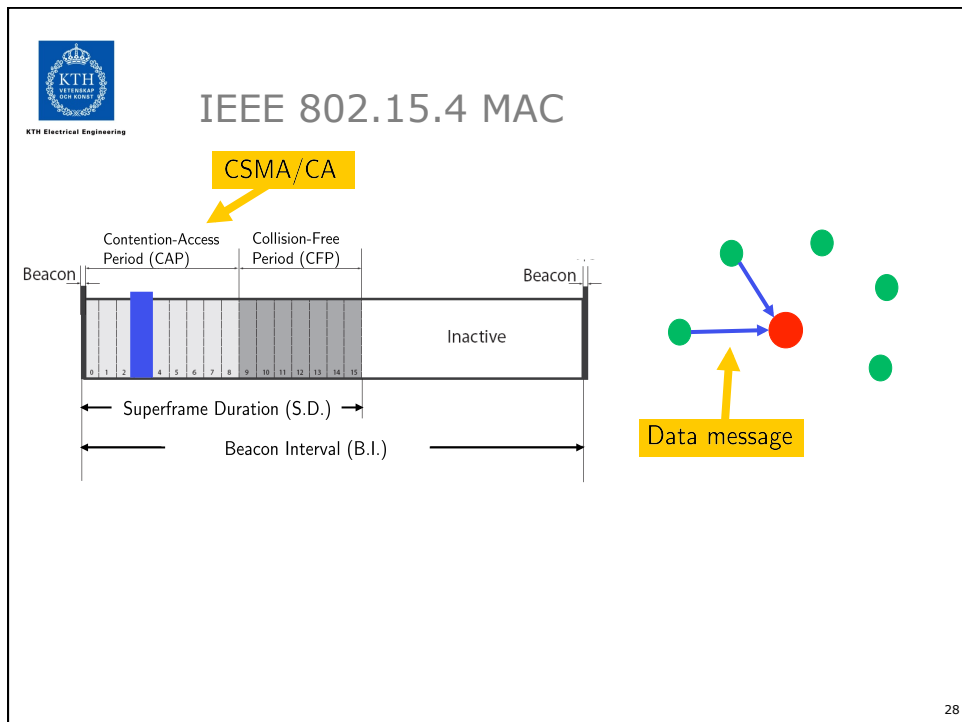
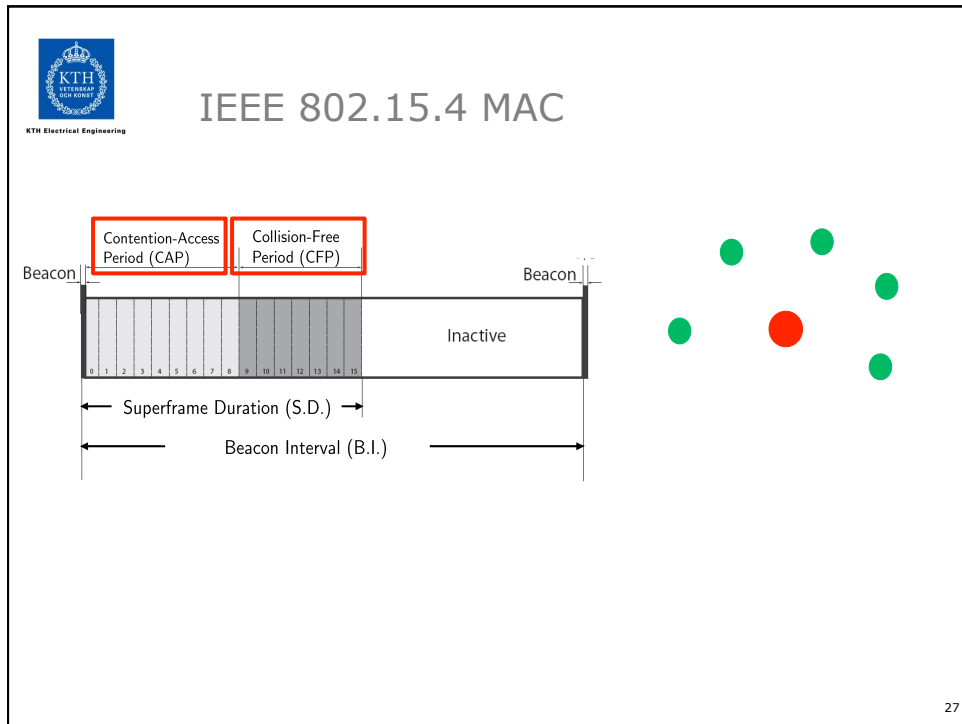
Contention access period for random CSMA communication

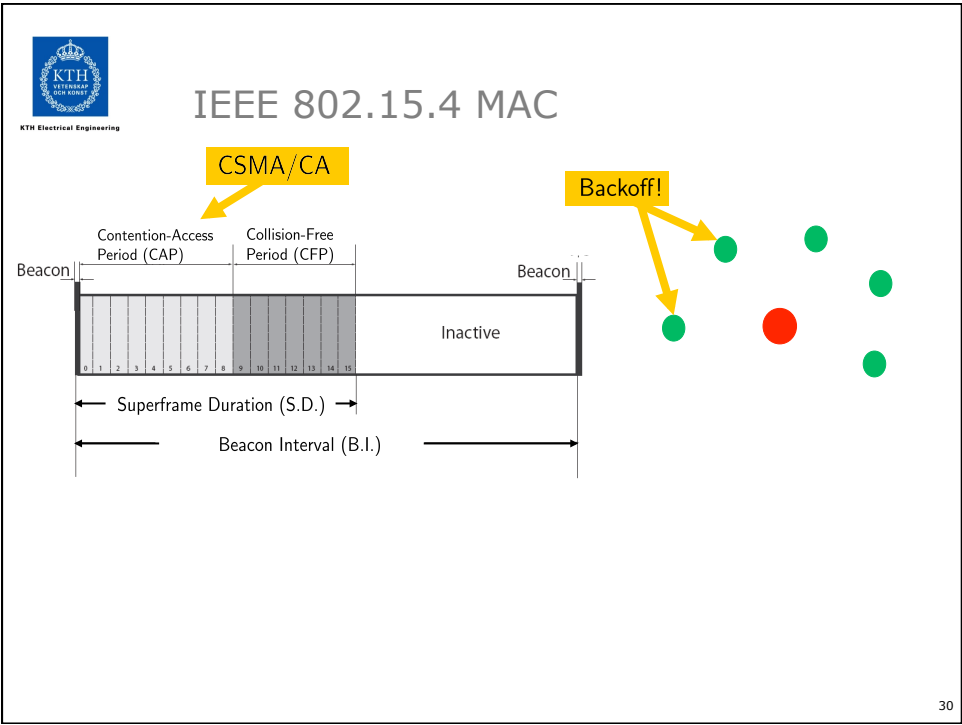
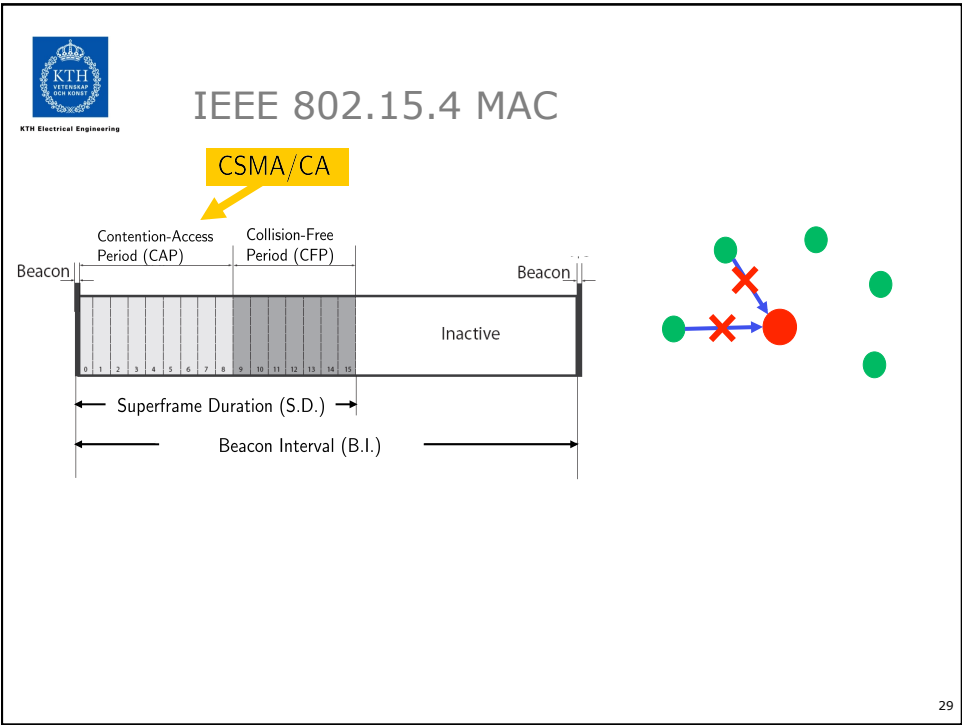


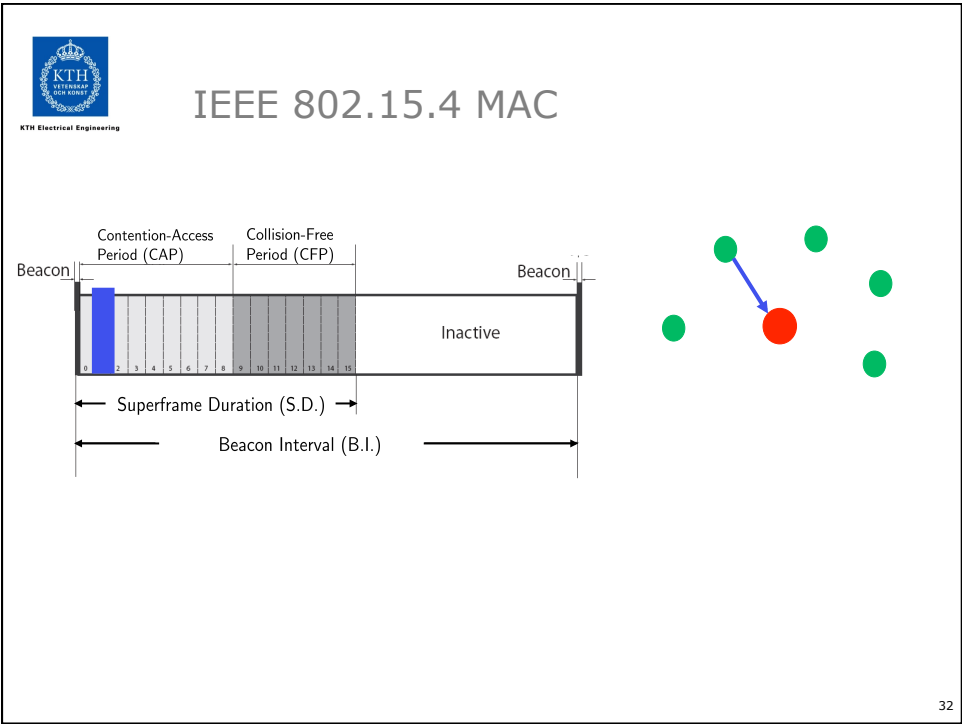
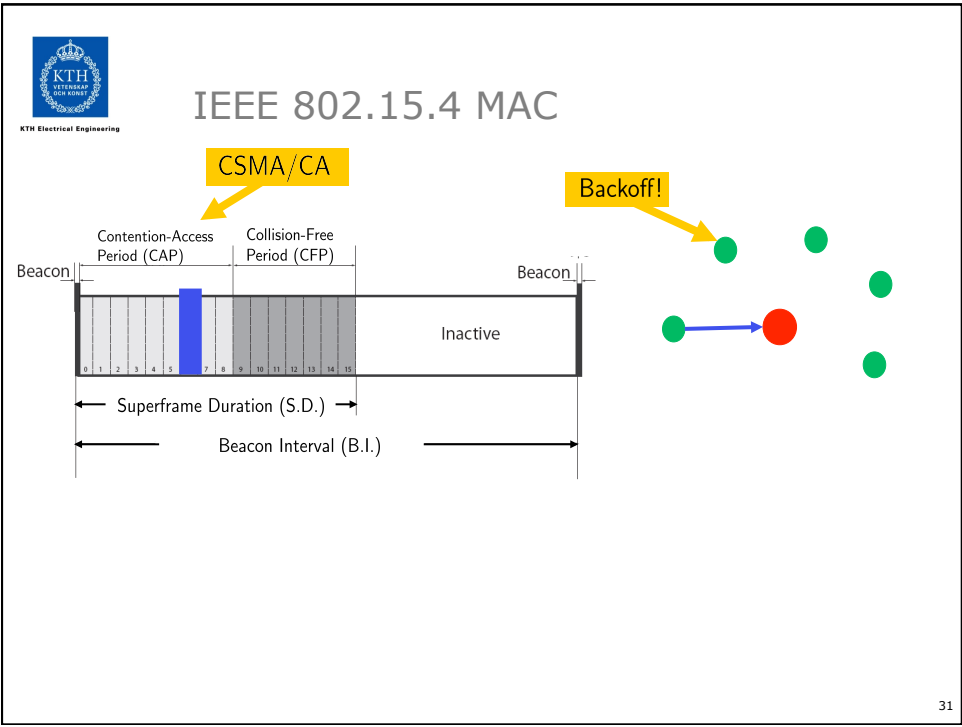
Araujo et al., 2010, Tiberi et al., 2010


TDMA = Time division multiple access, CSMA/CA = Carrier Sense Multiple Access with Collision Avoidance



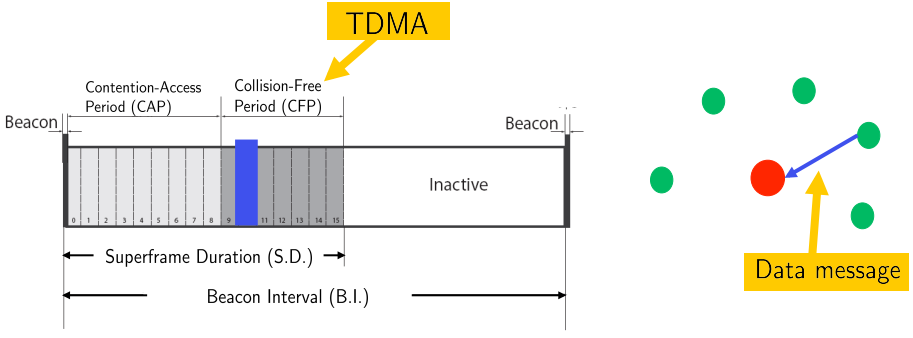







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IEEE 802.15.4 MAC

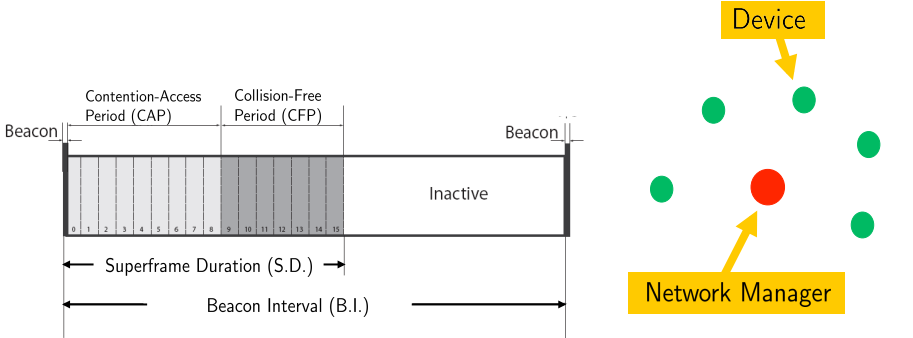


- CFP slot allocation as First-Come First-Served

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IEEE 802.15.4 MAC



- 16 slots for CAP and CFP
- Maximum 7 slots for CFP
- CFP slot allocation as First-Come First-Served

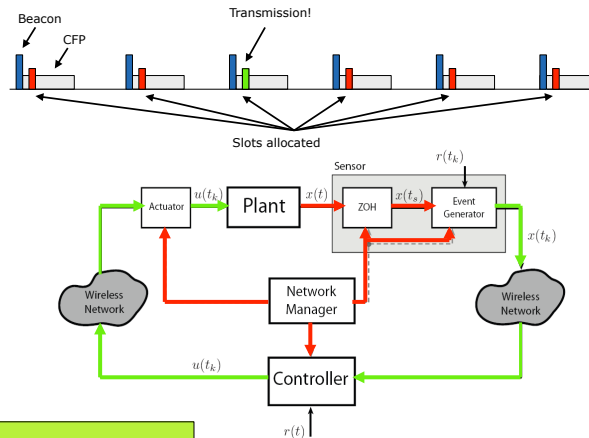
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Event-based sensor communication

1. Fixed scheduling of sensing/actuation slots
2. Check triggering condition at every allocated slot
 - One-step ahead triggering condition
3. If triggering condition is true, transmit measurement and perform actuation



- Robust to disturbances

- Unnecessary bandwidth utilization
- Energy spent on checking triggering condition

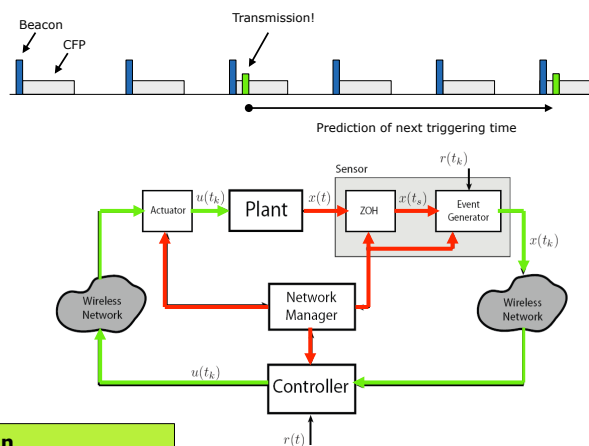
Araujo, 2011 35



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Predictive sensor communication

1. Scheduling of sensing/actuation slots when required, at beacon times
2. If triggering condition is predicted to be true, transmit measurement and perform control action
3. At every transmission, predict and schedule the next triggering time
 - Set node to sleep until next transmission



- Efficient bandwidth utilization
- Low energy consumption

- Less robust to disturbances

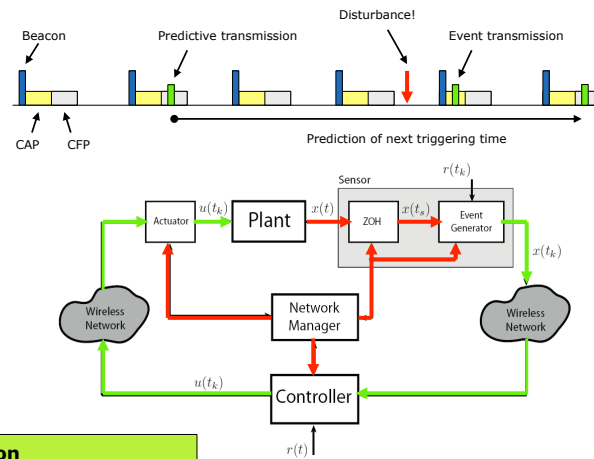
Araujo, 2011 36



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Hybrid sensor communication

1. Scheduling of slots as predictive scheme
2. Sensor node checks triggering condition continuously (or during CAP)
3. If triggering condition is true, transmit measurement and perform control action



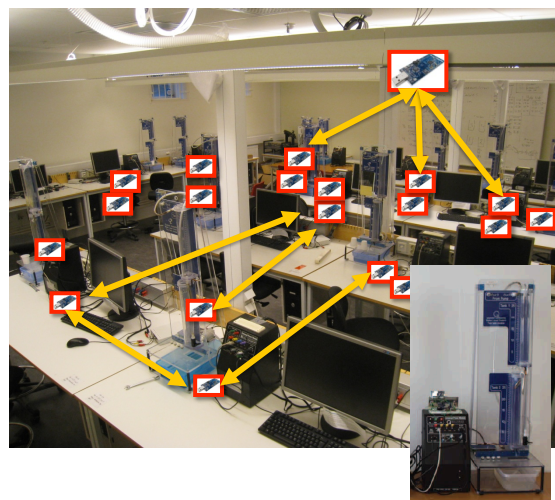
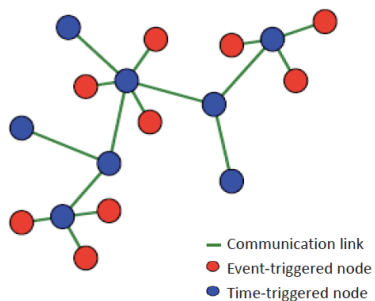
- Efficient bandwidth utilization
- Robust to disturbances

- Energy spent on checking triggering condition

Araujo, 2011 37

Multi-hop networks

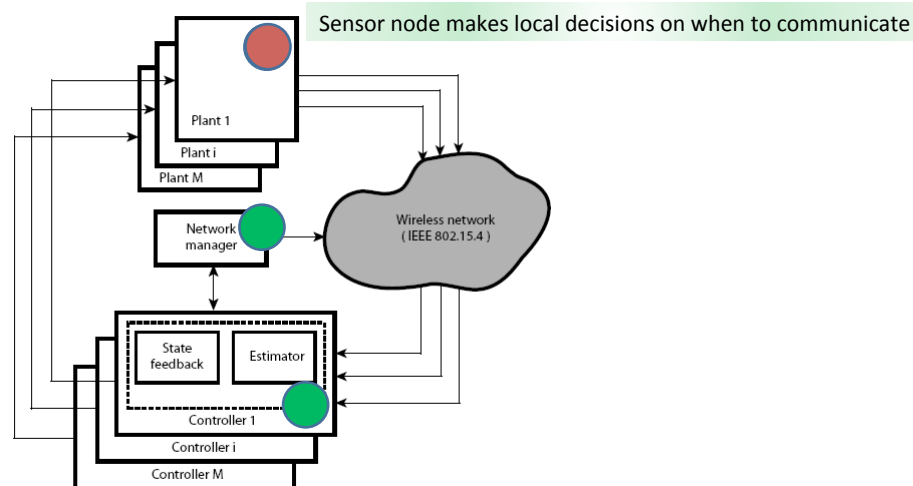
- Routing decisions
- Time delays
- Hidden terminal problem



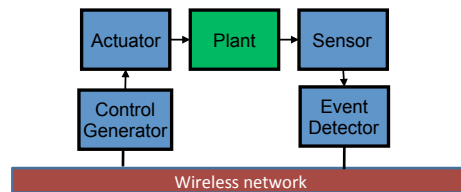
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Event-based control over wireless network



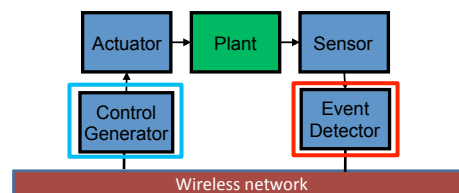
Event-based control loop



Åström, 2007, Rabi and J., WICON, 2008

When to transmit?

- Event detector mechanism on sensor side
 - E.g., threshold crossing



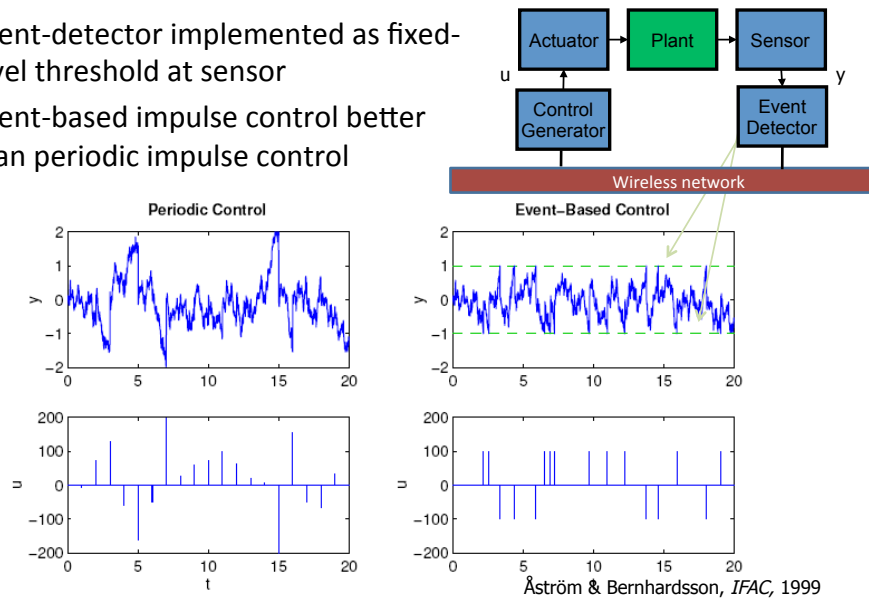
How to control?

- Execute control law at actuator side
 - E.g., piecewise constant controls, impulse control

Rabi et al., 2008

Example: Fixed threshold with impulse control

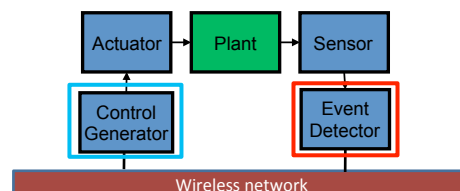
- Event-detector implemented as fixed-level threshold at sensor
- Event-based impulse control better than periodic impulse control



Control generators and event detectors

1. Impulse
2. Zero order hold
3. Higher order hold

1. Fixed threshold
2. Time-varying
3. Adaptive



Plant model

Plant

$$dx = udt + dv,$$

Stochastic differential equation, interpreted as

$$x(s + \tau) - x(\tau) = \int_{\tau}^{s+\tau} u(t)dt + \int_{\tau}^{s+\tau} dv(t)$$

with one ordinary (Lebesgue) integral and one stochastic (Ito) integral.

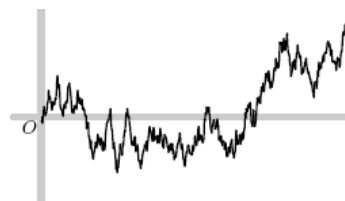
v is a Wiener process (or Brownian motion)

See Øksendal (2003) for an introduction to stochastic differential equations

Wiener process

A Wiener process $v(t)$ fulfills

1. $v(0)=0$
2. $v(t)$ is almost surely continuous
3. $v(t)$ has independent increments
with $v(t)-v(s) \sim N(0,t-s)$ for $t>s\geq 0$



Remark The variance of a Wiener process is growing like

$$E(v(t+s) - v(t))^2 = |s|$$

Plant model

Plant $dx = udt + dv,$

Stochastic differential equation, interpreted as

$$x(s + \tau) - x(\tau) = \int_{\tau}^{s+\tau} u(t)dt + \int_{\tau}^{s+\tau} dv(t)$$

with one ordinary (Lebesgue) integral and one stochastic (Ito) integral.

When $s > 0$ is a small, the change of $x(\tau)$ is normally distributed with mean $su(\tau)$ and variance s .

Plant model and control cost

Plant $dx = udt + dv,$

v is a Wiener process: $E(v(t + s) - v(t))^2 = |s|$

Cost function $V = \frac{1}{T} E \int_0^T x^2(t) dt.$

Periodic impulse control

Impulse applied at events t_k

$$u(t) = -x(t_k)\delta(t - t_k),$$

Periodic reset of state every event.

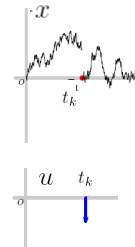
State grows linearly as

$$E(v(t+s) - v(t))^2 = |s|$$

between sample instances, because $dx = udt + dv$,

Average variance over sampling period h is $\frac{1}{2}h$ so the cost is

$$V_{PIH} = \frac{1}{2}h.$$



Åström, 2007

Periodic ZoH control

Traditional sampled-data control theory gives that

$V = \frac{1}{h} \int_0^h E x^2(t) dt$ is minimized for the sampled system

$$x(t+h) = x(t) + hu(t) + e(t),$$

with

$$u = -Lx = \frac{1}{h} \frac{3 + \sqrt{3}}{2 + \sqrt{3}} x$$

derived from

$$S = \Phi^T S \Phi + Q_1 - L^T R L, \quad L = R^{-1}(\Gamma^T S \Phi + Q_{12}^T), \quad R = Q_2 + \Gamma^T S \Gamma,$$

The minimum gives the cost

$$V_{PZO H} = \frac{3 + \sqrt{3}}{6} h$$

Åström, 2007

Event-based impulse control with fixed threshold

Suppose an event is generated whenever

$$|x(t_k)| = a$$

generating impulse control

$$u(t) = -x(t_k)\delta(t - t_k),$$

One can show that the average time
between two events is

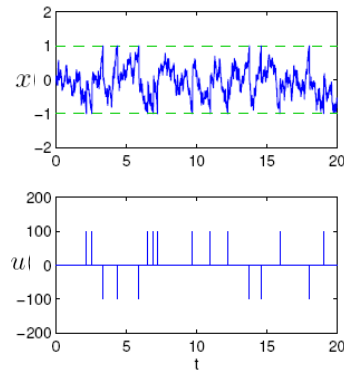
$$h_E := E(T_{\pm d}) = E(x_{T_{\pm d}}^2) = a^2$$

and that the pdf of x is triangular:

$$f(x) = (a - |x|)/a^2$$

The cost is

$$V_{EIH} = \frac{a^2}{6} = \frac{h_E}{6}$$

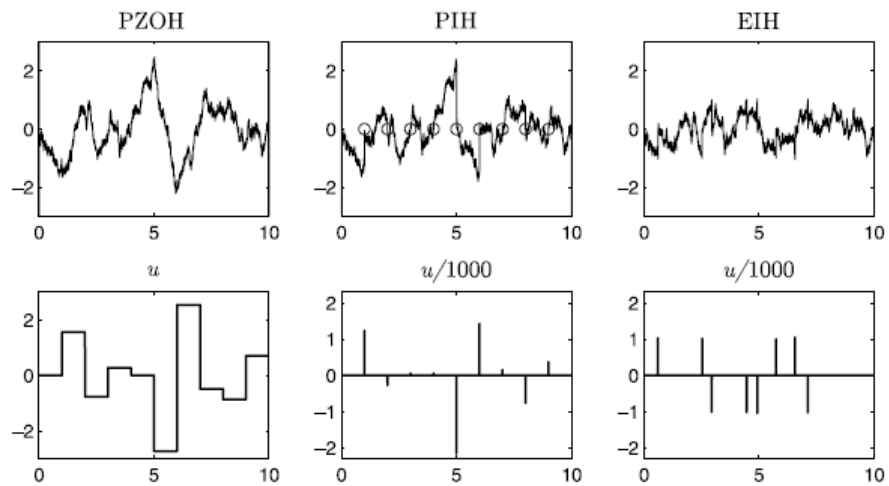


Åström, 2007

Pdf $f(x) = (a - |x|)/a^2$ is the solution to the forward
Kolmogorov forward equation (or Fokker–Planck
equation)

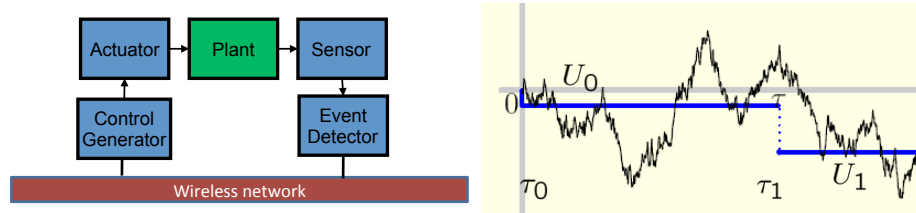
$$\frac{\partial f}{\partial t} = \frac{1}{2} \frac{\partial^2 f}{\partial x^2}(x) - \frac{1}{2} \frac{\partial f}{\partial x}(d)\delta_x + \frac{1}{2} \frac{\partial f}{\partial x}(-d)\delta_x, \quad f(-a) = f(a) = 0,$$

Comparison



Åström, 2007

Event-based ZoH control with adaptive sampling



How choose $\{U_i\}$ and $\{\tau_i\}$ to minimize $V = \frac{1}{T} E \int_0^T x^2(t) dt$.

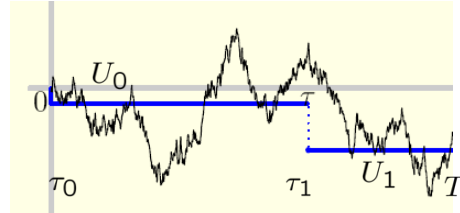
Rabi et al., 2008

Optimal control with one sampling event

$$dx_t = u_t dt + dB_t$$

$$\min_{U_0, U_1, \tau} J = \min_{U_0, U_1, \tau} \mathbf{E} \int_0^T x_s^2 ds$$

$$= \min_{U_0, U_1, \tau} \left[\mathbf{E} \int_0^\tau x_s^2 ds + \mathbf{E} \int_\tau^T x_s^2 ds \right]$$

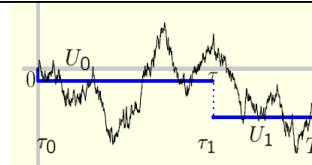


A joint optimal control and optimal stopping problem

Rabi et al., 2008

$$dx_t = u_t dt + dB_t$$

$$\min_{U_0, U_1, \tau} J = \min_{U_0, U_1, \tau} \mathbf{E} \int_0^T x_s^2 ds$$



If τ chosen deterministically (not depending on x_t)
and $x_0 = 0$:

$$U_0^* = 0 \quad U_1^* = -\frac{3x_{T/2}}{T} \quad \tau^* = T/2$$

If τ is event-driven (depending on x_t) and $x_0 = 0$:

$$U_0^* = 0 \quad U_1^* = -\frac{3x_{\tau^*}}{2(T - \tau^*)}$$

$$\tau^* = \inf \{t : \underbrace{x_t^2}_{\geq \sqrt{3}(T-t)} \geq \sqrt{3}(T-t)\}$$

Rabi et al., 2008

Envelope defines optimal level detector

The diagram illustrates a control system and a corresponding stochastic process plot. The control system consists of an Actuator, Plant, Sensor, and Level Detector, all connected via a Wireless network. The Level Detector is also connected to a Dynamic level detector. The plot shows a stochastic process x_t over time t , with a horizontal line representing the level u . The process is bounded by a parabolic curve. The initial condition is $x_0 = 0$ and the final condition is $x_T = 0$. The process is defined by the stochastic differential equation $dx_t = u_t dt + dB_t$ and the cost function $J = \min_{U_0, U_1, \tau} \mathbf{E} \int_0^T x_s^2 ds$.

Multiple samples

$$J_N \left(x_0, \mathcal{U}, \{\tau\}_{i=1}^N \right) = \mathbb{E} \left[\int_0^T x_s^2 ds \middle| x_0 \right]$$

Extension to variable budget sampling, allowing number of samples to depend on x .

Event-based impulse control over wireless network with communication losses

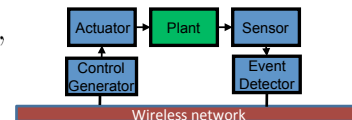
Plant $dx_t = dW_t + u_t dt, x(0) = x_0,$

Sampling events $\mathcal{T} = \{\tau_0, \tau_1, \tau_2, \dots\},$

Impulse control $u_t = \sum_{n=0}^{\infty} x_{\tau_n} \delta(\tau_n)$

Average sampling rate $R_\tau = \limsup_{M \rightarrow \infty} \frac{1}{M} \mathbb{E} \left[\int_0^M \sum_{n=0}^{\infty} \mathbf{1}_{\{\tau_n \leq M\}} \delta(s - \tau_n) ds \right]$

Average cost $J = \limsup_{M \rightarrow \infty} \frac{1}{M} \mathbb{E} \left[\int_0^M x_s^2 ds \right]$



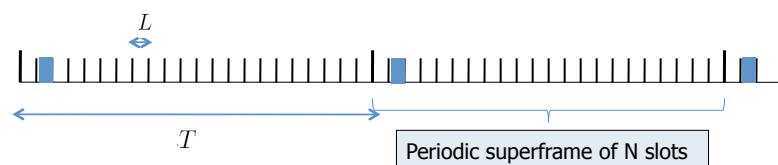
Periodic impulse control

Sampling events $\tau_n = nT$ for $n \geq 0$

Slot length L gives $T = NL$

Average sampling rate $R_{\text{Periodic}} = \frac{1}{T}$

Average cost $J_{\text{Periodic}} = \frac{T}{2}$

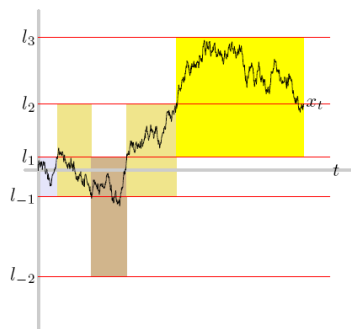


Level-triggered event-based control

Ordered set of levels $\mathcal{L} = \{\dots, l_{-2}, l_{-1}, l_0, l_1, l_2, \dots\}$ $l_0 = 0$

Multiple levels needed because we allow packet loss

Lebesgue sampling $\tau = \inf \{ \tau > \tau_i, x_\tau \in \mathcal{L}, x_\tau \notin x_{\tau_i} \}$



Level-triggered control

For Brownian motion, equidistant sampling is optimal

$$\mathcal{L}^* = \{k\Delta \mid k \in \mathbb{Z}\}$$

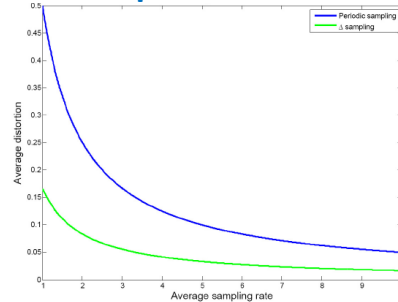
First exit time

$$\tau_\Delta = \inf \{ \tau \geq 0, x_\tau \notin (\xi - \Delta, \xi + \Delta), x_0 = \xi \}$$

Average sampling rate $R_\Delta = \frac{1}{\mathbb{E}[\tau_\Delta]} = \frac{1}{\Delta^2},$

Average cost $J_\Delta = \frac{\mathbb{E}[\int_0^{\tau_\Delta} x_s^2 ds]}{\mathbb{E}[\tau_\Delta]} = \frac{\Delta^2}{6}.$

Comparison between **periodic** and **event-based** control



$T = \Delta^2$ gives equal average sampling rate for periodic control and event-based control

Event-based impulse control is 3 times better than periodic impulse control

What about the influence of communication losses?
When is event-based sampling better and vice versa?

Influence of communication losses

Times when packets are successfully received $\rho_i \in \{\tau_0 = 0, \tau_1, \tau_2, \dots\}$,

$$\{\rho_0 = 0, \rho_1, \rho_2, \dots\}, \quad \rho_i \geq \tau_i,$$

Average rate of packet reception

$$R_p = \limsup_{M \rightarrow \infty} \frac{1}{M} \mathbb{E} \left[\int_0^M \sum_{n=0}^{\infty} \mathbf{1}_{\{\rho_n \leq M\}} \delta(s - \rho_n) ds \right] = p \cdot R_\tau$$

Define the times between successful packet receptions $\rho_{(p, \Delta)}$

$$\text{Average cost } J_p = \limsup_{T \rightarrow \infty} \frac{1}{T} \mathbb{E} \left[\int_0^T x_s^2 ds \right] = \frac{\mathbb{E} \left[\int_0^{\rho_{(p, \Delta)}} x_s^2 ds \right]}{\mathbb{E} [\rho_{(p, \Delta)}]}$$

Periodic control under packet losses

Sampling with fixed period T with loss probability p gives cost

$$J = \frac{T(1+p)}{2(1-p)}$$

Compared with event-based control by setting

$$T = \Delta^2$$

so that the average use of the communication channel is equal

Event-based control under packet losses

Proposition

If packet losses are IID with prob p , then equidistant event-based (Lebesgue) sampling gives

$$J_p = \frac{\Delta^2 (5p + 1)}{6(1-p)}$$

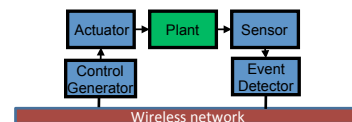
Remark

- Event-based control with losses always better than periodic with losses.
- Event-based control with losses outperformed by periodic control without losses if

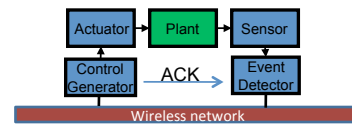
$$\frac{(1+5p)}{3(1-p)} \geq 1$$

so if $p \geq 0.25$ then periodic sampling do better than event-based sampling.

Rabi and J., 2009



Sensor data ACK's



If controller perfectly acknowledges packets to sensor,
event detector can adjust its sampling strategy

Let $\Delta(l) = \sqrt{l+1}\Delta_0$

where $l \geq 0$ number of samples lost since last successfully
transmitted packet

Gives $\mathbb{E}[\tau_{i+1}^\uparrow - \tau_i^\uparrow]$ independent of i .

Better performance than fixed $\Delta(l)$ for same sampling rate:

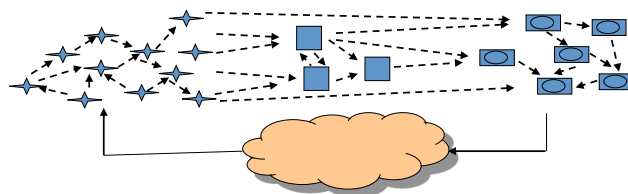
$$J_p^\uparrow = \frac{\Delta^2(1+p)}{6(1-p)} \leq \frac{\Delta^2(1+5p)}{6(1-p)} = J_p.$$

Lecture 2 Outline

- What's new with wireless networked control?
- State-based scheduling for control
- Exploiting wireless protocols for control
- Event-based control
- Conclusions

Conclusions

- **Wireless control and networking** are enabling technologies in many emerging industrial applications
- Fundamental challenges related to
 - **time-driven**, synchronous, sampled data control theory, vs
 - **event-driven**, asynchronous, ad hoc wireless networking
- New principles for control in large-scale wireless systems



Take-home message

Lecture 1: Motivating applications and challenges

- Networked control systems have societal importance
- Many new applications with challenging problems

Lecture 2: Wireless control systems

- Everything will be wireless, including control systems
- Interesting research challenges on the intersection between sensor networks, wireless communication, and control theory

<http://www.ee.kth.se/~kallej>

Outline

Lecture 1: Motivating applications and challenges

Lecture 2: Wireless control systems

Take-home message

Lecture 1: Motivating applications and challenges

- Networked control systems have societal importance
- Many new applications with challenging problems

Lecture 2: Wireless control systems

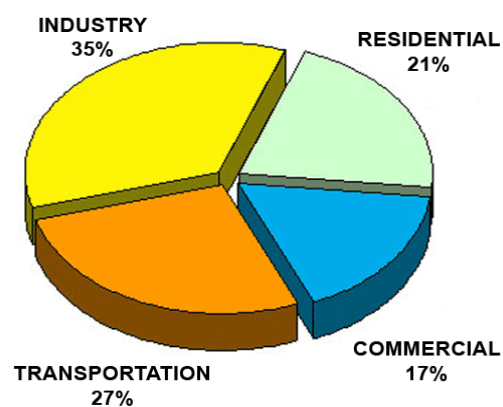
- Everything will be wireless, including control systems
- Interesting research challenges on the intersection between sensor networks, wireless communication, and control theory

Outline

Lecture 1: Motivating applications and challenges

Lecture 2: Wireless control systems

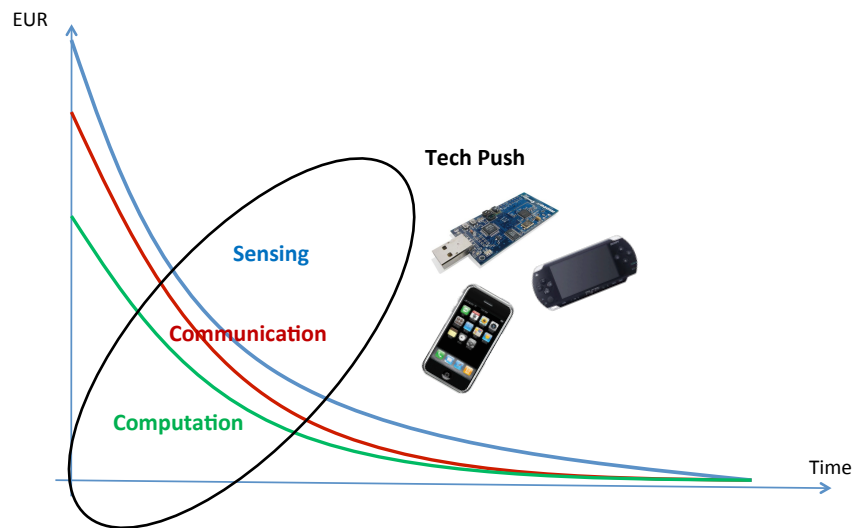
Energy consumption



More and better networked control reduces energy consumption

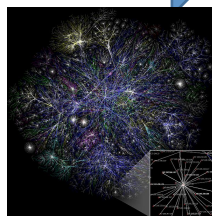
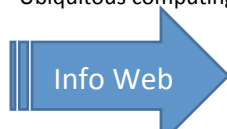
US Energy Information Administration

Why now?



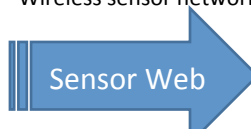
Network evolution

- Internet
- WWW
- Ubiquitous computing



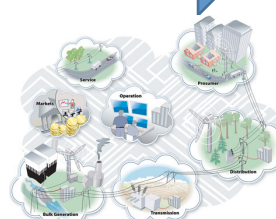
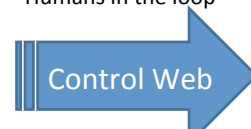
The Internet

- Remote sensing
- Monitoring environments
- Wireless sensor networks

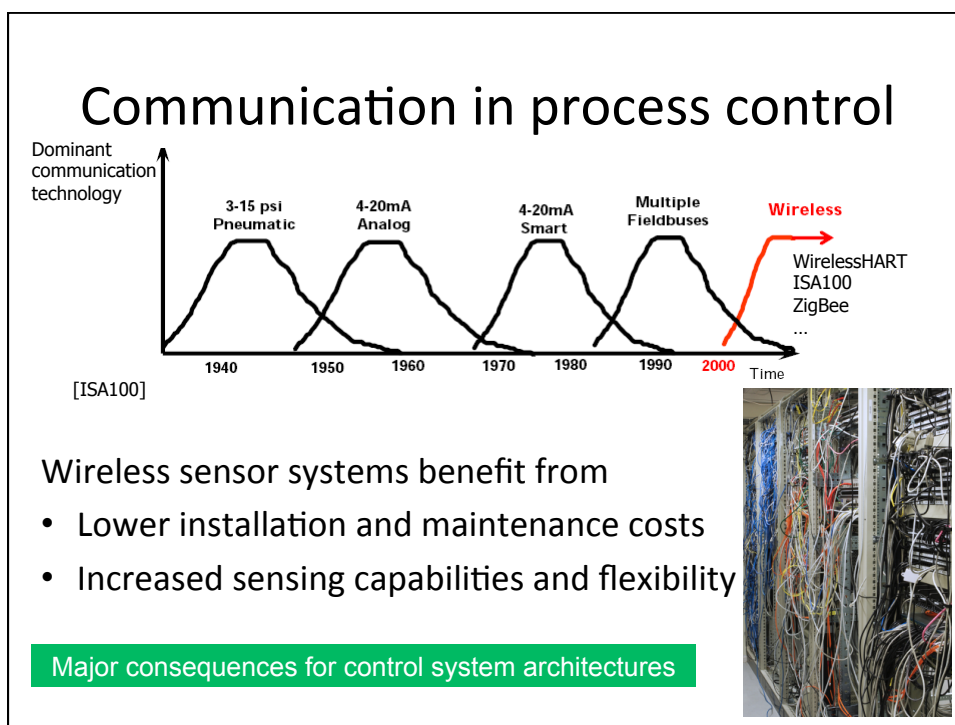
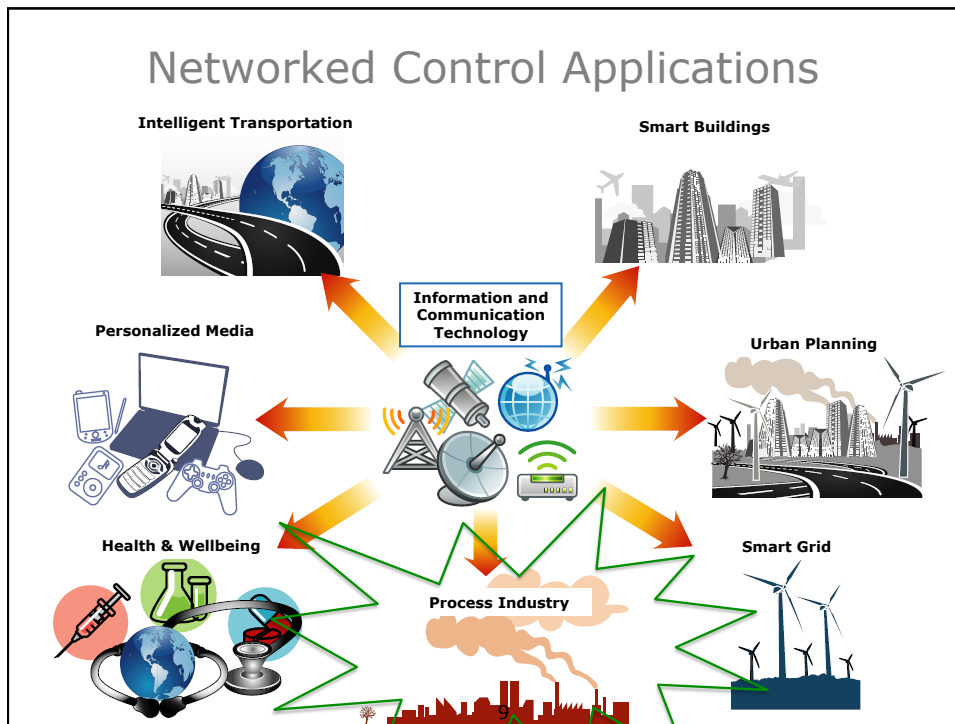


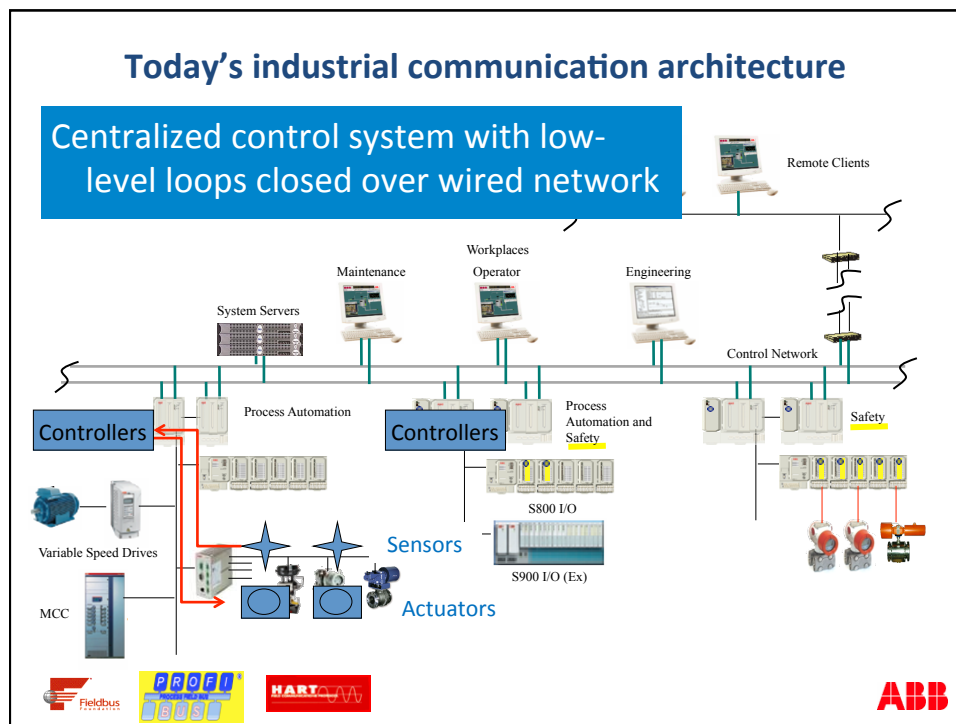
Monitoring storm petrels at Great Duck Island

- Closing the loop
- Cyber-physical systems
- Humans in the loop



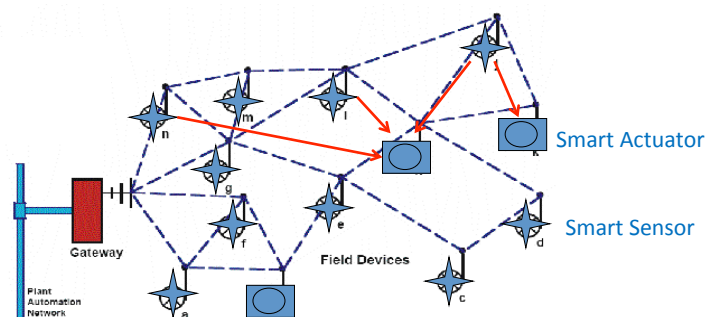
The smart energy grid





Towards wireless sensor and actuator network architecture

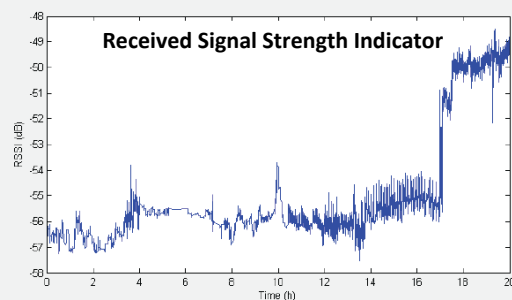
- Local control loops closed over **wireless multi-hop network**
- Potential for a dramatic change:
 - From fixed hierarchical centralized system to flexible distributed
 - Move intelligence from dedicated computers to sensors/actuators



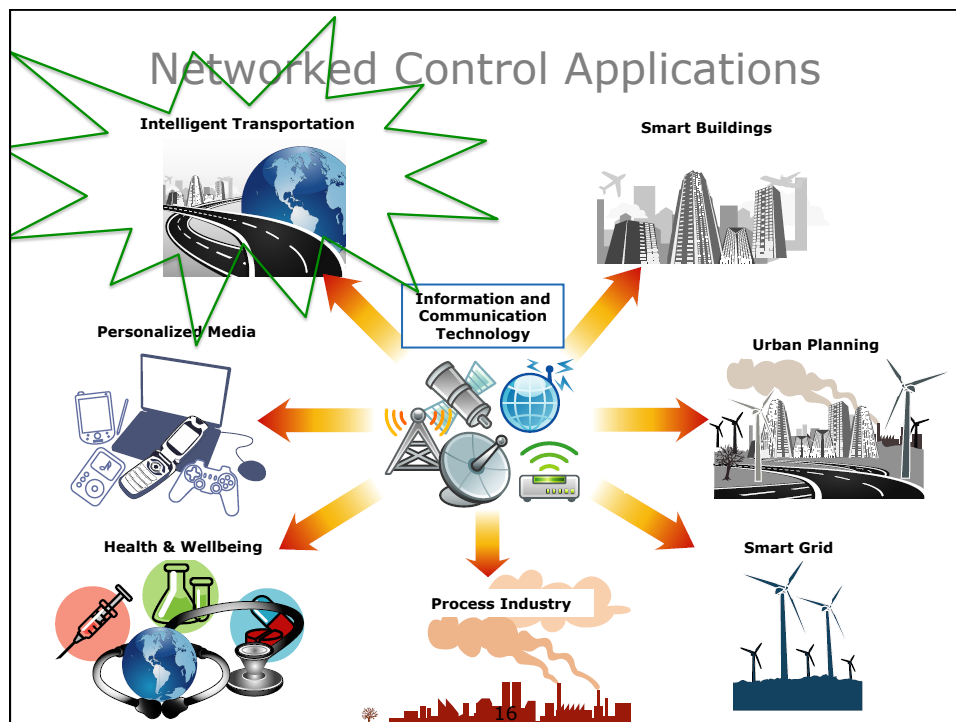
Radio Channel Measurements in Industrial Environment

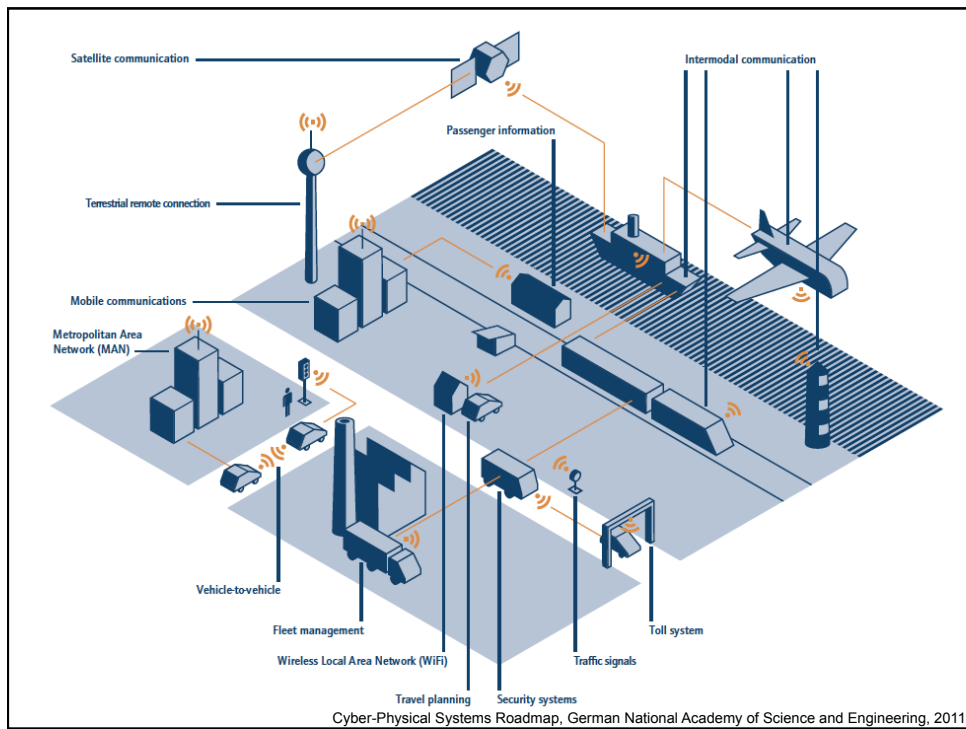


- Rolling mill at Sandvik in Sweden
- Study of 2.45 GHz radio channel properties
- Slow but substantial RSSI variations due to mobile machines



Ahlen et al, 2012



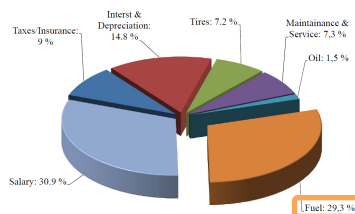


Demands from Goods Road Transportation

- Goods transportation accounts for 30% of CO₂ emissions
15% of greenhouse gas emissions of the global fossil fuel combustion
- Expected to increase by 50% for 2000-2020

International Transport Forum (2010), EC (2006)

Life cycle cost for European heavy-duty vehicles



Total fuel cost 80 k€/year/vehicle

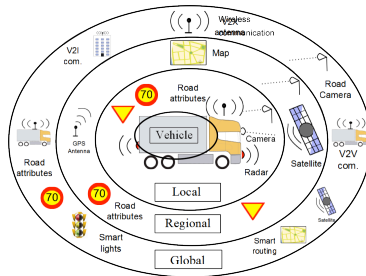
Schittler, 2003

24% of long haulage trucks run empty
57% average load capacity

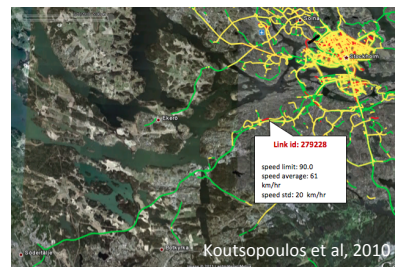
Dr. H. Ludanek, CTO, Scania

Technology Push

Sensor and communication technology



Real-time traffic information

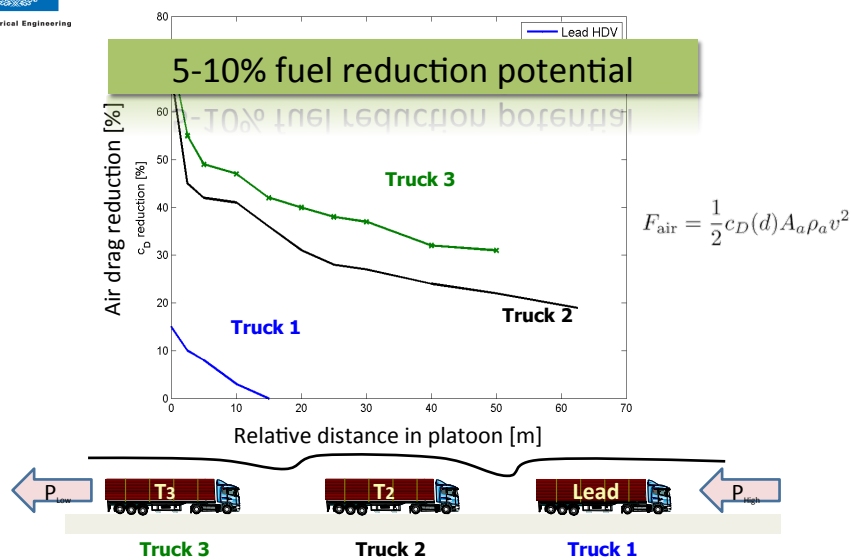


Vehicle platooning and semi-autonomous driving



Air Drag Reduction in Platooning

5-10% fuel reduction potential

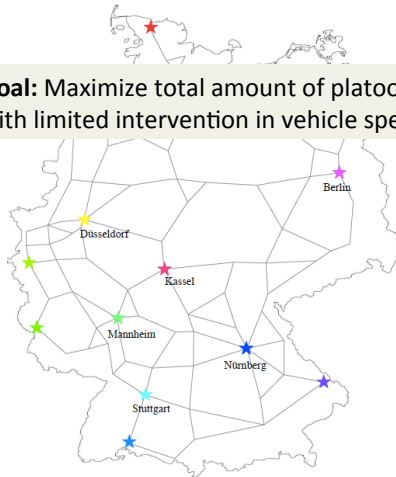


Wolf-Heinrich & Ahmed (1998), Bonnet & Fritz (2000), Scania CV AB (2011)

Fuel-Optimal Goods Transportation

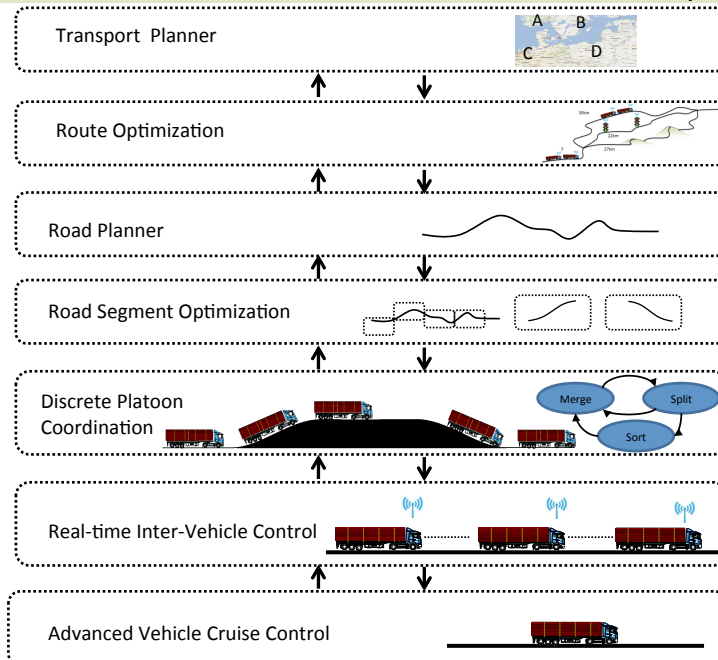
- Goods transported between cities over European highway network
- 2 000 000 long haulage trucks in EU (400 000 in Germany)
- Large distributed control systems with no real-time coordination today

Goal: Maximize total amount of platooning with limited intervention in vehicle speed and route



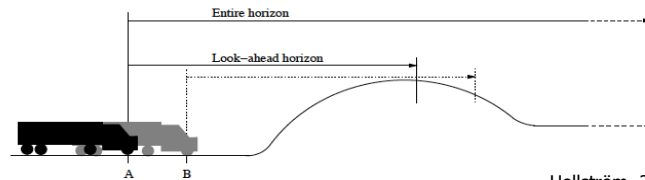
Larson et al., 2013

Architecture for Future Coordinated Goods Transportation



Alam et al., 2012

Look-ahead Cruise Control for Individual Vehicle



Hellström, 2007

Adjust driving force to **minimize fuel consumption based on road topology** info:

The total fuel consumption over time T is:

$$\int_0^T \delta(t) \left(\frac{1}{\eta} \frac{dv(t)}{dt} + \frac{1}{2} \rho_a A_a c_D v^2(t) + mgc_r \cos \alpha + mg \sin \alpha \right) dt \quad (3)$$

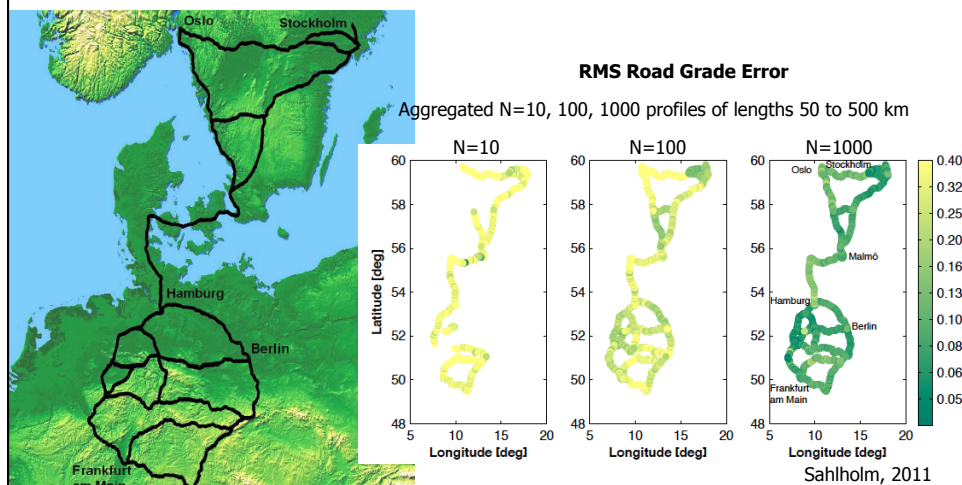
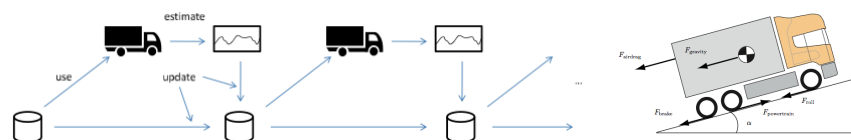
Require knowledge of road grade α , not available in today's navigators

$$\begin{aligned} m_t \frac{dv}{dt} &= F_{eng} - F_b - F_{ad}(v, d) - F_r(\alpha) - F_g(\alpha) \\ &= F_{eng} - F_b - \frac{1}{2} \rho_a A_a c_D v^2 \phi(d) \\ &\quad - mgc_r \cos \alpha - mg \sin \alpha \end{aligned}$$

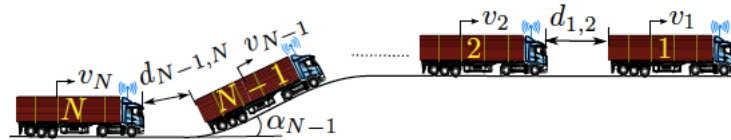
Implemented as reference change to existing cruise controller

Allam et al., 2011

Distributed Road Grade Estimation



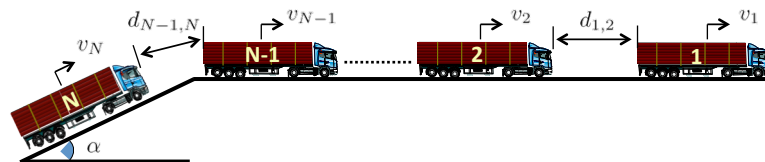
Look Ahead Cruise Control for Platoons



- How to jointly minimize fuel consumption for a platoon of vehicles?
 - Uphill and downhill segments
- How to order vehicles according to weight and other performance criteria?
 - Heavy and light vehicles

Alam et al., 2013

How to Control Vehicles in a Platoon?



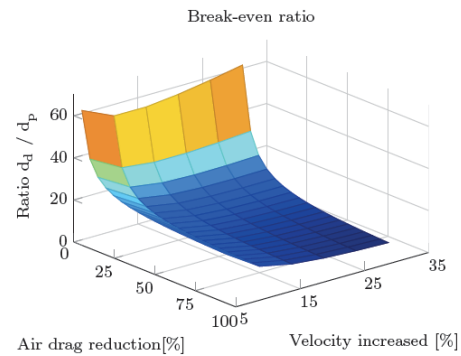
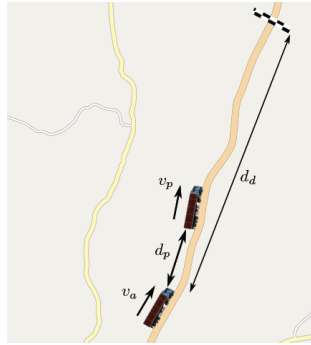
- Platooning control applications require **collaborative actions**
- Vehicles need **accurate estimates** of neighboring vehicles' states and actions
 - Vehicle-to-vehicle and vehicle-to-infrastructure communications
 - Wireless communication standards, e.g., IEEE 802.11p
- Control performance is tightly coupled to how well state information (position, velocity, braking etc) is communicated across the platoon



- How does the communication influence the system performance?
- What is an efficient communication strategy for specific control tasks?

Extensive theoretical and practical studies since Levine & Athans, 1966

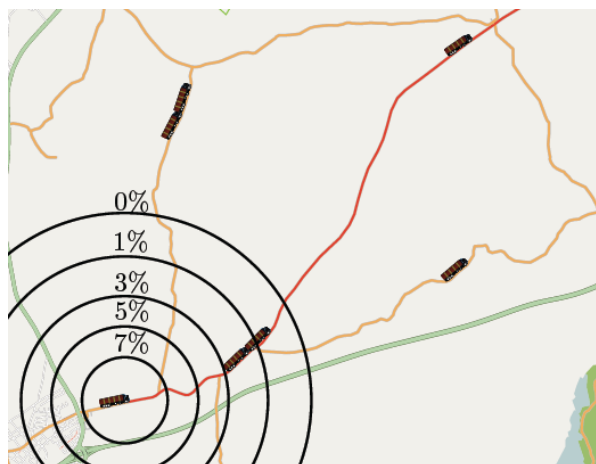
When is it Fuel Efficient for a Heavy-Duty Vehicle to Catch Up with a Platoon?



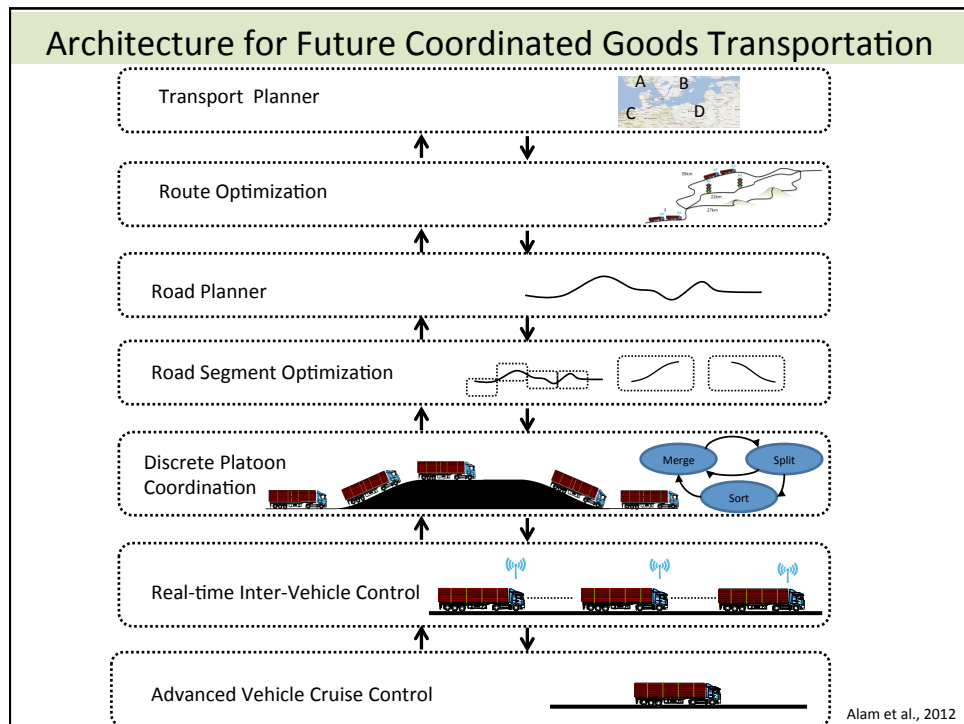
- Catch up costs fuel due to higher air drag at higher velocity
- Distance to platoon d_p needs to be small compared to total travel distance d_d
- Tradeoff velocity increase during catch-up and air drag reduction during platooning

Liang et al., 2013

When is it Fuel Efficient for a Heavy-Duty Vehicle to Catch Up with a Platoon?



Liang et al., 2013



When create platoons?



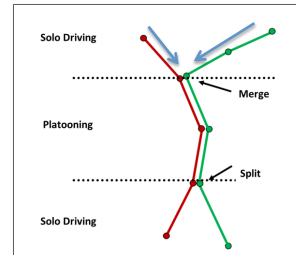
Larson et al., 2013

Platoon merge and split

Heavy-duty vehicle traffic without platooning



Merge and split platoons at highway intersections

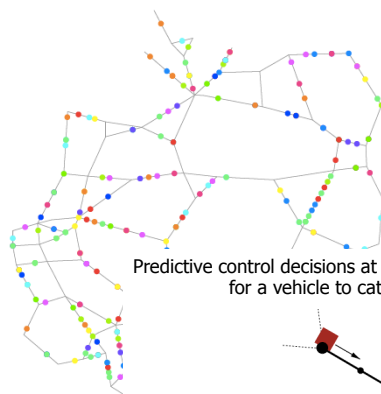


Only vehicles that are relatively close in space and time platoon

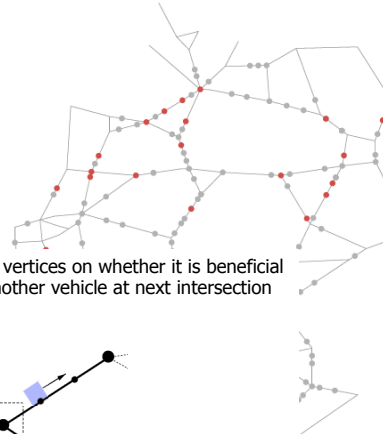
Larson et al., 2013

Distributed optimization of platooning

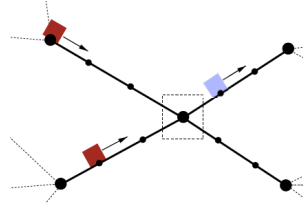
Heavy-duty vehicle traffic without platooning



With platooning

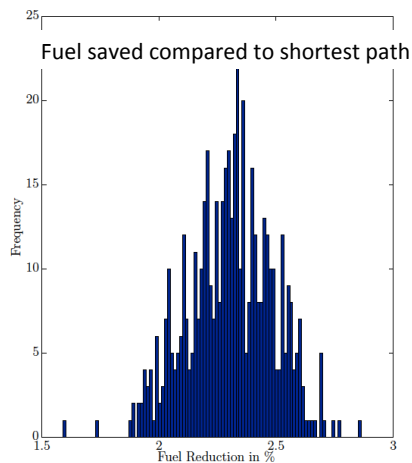


Predictive control decisions at network vertices on whether it is beneficial for a vehicle to catch up another vehicle at next intersection

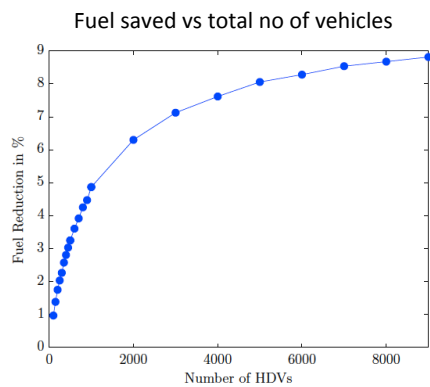


Larson et al., 2013

Numerical evaluations



- German road network with 300 trucks
- Random starting points and destinations
- 500 Monte Carlo experiments



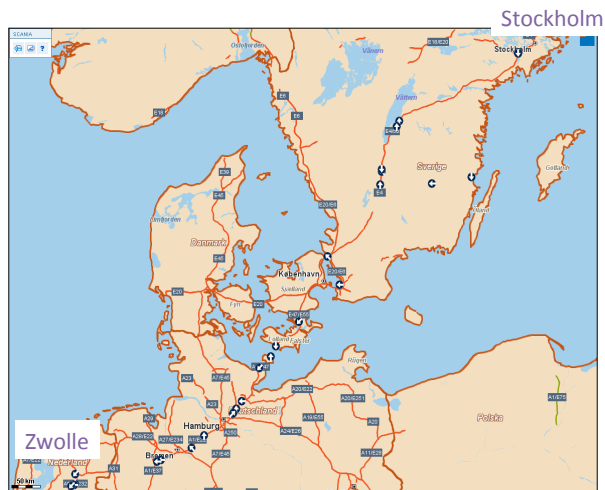
2-5% deployment enough for substantial benefit

Scania Transport Lab 2013

Stockholm-Zwolle Testsite

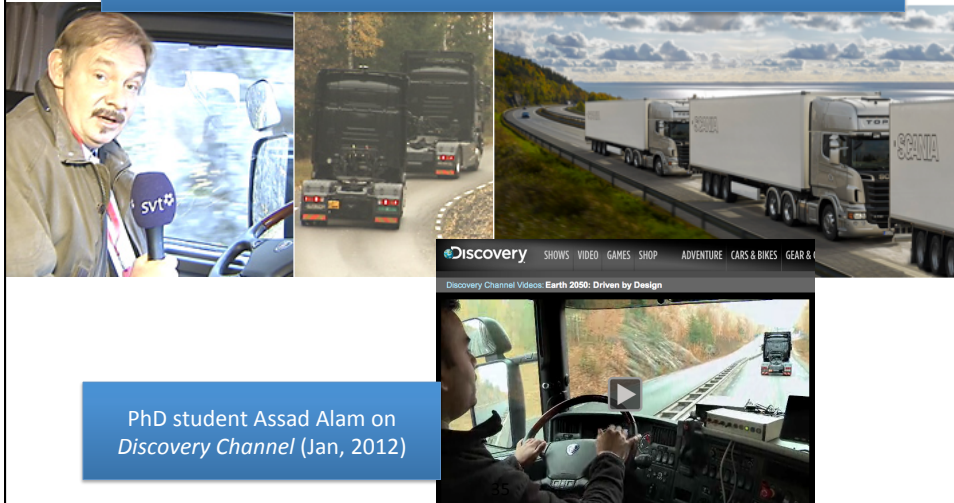
- Real-time fleet management
- Platooning in real traffic
- Fuel reductions and safety
- Driver acceptance
- Public acceptance

Scania Transport Lab
Internal haulage company
20 trucks, 360.000 km/year
75 trailers, 92% loaded
65 drivers, 40 h work/week



Platooning Demos in Media

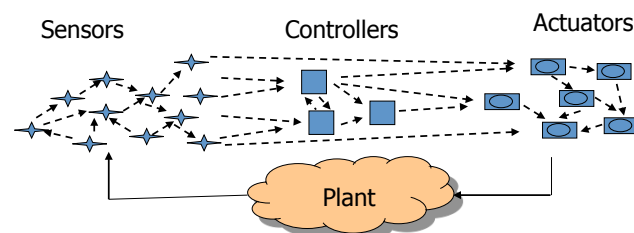
Report on vehicle platooning developed by KTH and Scania (Oct, 2011)



PhD student Assad Alam on
Discovery Channel (Jan, 2012)

Wireless control system

How share common network resources while
maintaining guaranteed closed-loop performance?



- How handle network imperfections: resource constraints, loss, conflicts, delays, outages?
- How move intelligence from a few central units to many distributed devices?